Robust Energy Efficiency Optimization for Amplify-and-Forward MIMO Relaying Systems

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Abstract-We investigate the energy efficiency (EE) of multiple-input-multiple-output (MIMO) amplify-and-forward relaying networks relying on the realistic imperfect channel state information (CSI). Specifically, the relay jointly optimizes the source covariance and relay beamforming matrices by maximizing the EE under additive or multiplicative relay-destination CSI errors. The optimal channel-diagonalizing structure is derived for the source covariance and relay beamforming matrices under the spectral-norm constrained additive or multiplicative CSI error. Then, the existence of a saddle point is proved, which shows that the channel-diagonalizing transmission strategy is optimal in the robust EE maximization under these two types of CSI errors, and the original matrix-valued fractional robust EE problem is transformed into a scalar fractional problem. We propose the Dinkelbach method-based alternating optimization scheme for this transformed robust EE problem, which is capable of finding a locally optimal solution of the original robust EE problem efficiently, and show that the semi-closedform solution to each of the two associated subproblems can be obtained. We then prove that the channel-diagonalizing transmission strategy remains optimal when the statistically imperfect source-relay channel is additionally imposed. We also extend our work into multi-hop MIMO relaying scenarios and prove that the channel-diagonalizing structure is optimal for the source covariance matrix and the multiple relay beamforming matrices.

Index Terms—Robust energy efficiency optimization, additive and multiplicative CSI errors, channel-diagonalization.

I. INTRODUCTION

COOPERATIVE relaying is a promising technique for improving the communication reliability and expanding

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the communication range [1]–[3]. Moreover, given the multiplexing and/or diversity gains provided by multiple-input multiple-output (MIMO) techniques, various relaying strategies have been proposed for MIMO relaying systems [3]-[8]. Among these existing relaying strategies, the amplify-andforward (AF) strategy is popular since only a simple linear transformation is required for forwarding signals by relays [3], [5]. Most of the existing literature of MIMO AF relaying systems concentrates on the optimization of traditional performance metrics, such as the achievable capacity and the minimum mean square errors (MSE) of signal detection [4]-[8]. Recently, considerable attention has been focused on the energy efficiency (EE), which is an important system performance metric for promoting green communications. Traditionally, it is defined as the ratio of the achievable capacity to the total power consumption of signal transmission and circuit hardware dissipation [9]-[12]. There exist some works in the literature that investigate the EE optimization for MIMO AF relaying systems [10]–[12]. Interestingly, these works provide a common insight that the channel-diagonalizing transceiver structure is optimal, in terms of EE optimization, which implies that similar to the optimization of the traditional capacity and MSE metrics [7], [8], the eigenmode transmission strategy is still optimal for EE optimization. However, these works are based on the unrealistic assumption of perfect channel state information (CSI). Since the CSI estimation errors are generally unavoidable in practice, ignoring this uncertainty, as in the works [10]–[12], will lead to significant EE performance degradation for MIMO systems. Consequently, it is necessary to consider the influence of CSI errors on the EE optimization of MIMO AF relaying systems.

There are two types of imperfect CSI models. One is the statistical CSI model, in which only partial CSI is available, such as channel mean or covariance matrix. In this context, given the channel distribution, various designs based on system average performance were investigated [11], [13], [14]. For example, the works [11], [13] studied the robust EE maximization of two-hop MIMO relaying networks given the statistical source-relay channel or relay-destination channel. In this case, the eigenmode transmission strategy is optimal. The other is the approximate CSI model, which adopts an error model for the CSI approximation. Generally, various CSI error models are classified into the deterministic and stochastic ones [15]–[21]. The deterministic CSI error is often used for modeling quantization inaccuracy, in which only

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the estimated channel knowledge with bounded CSI error is available [15], [16]. In this case, the worst-case robustness optimization is mainly considered [17]-[19]. In other words, the system performance under the worst-case channel quality becomes an important criterion. Naturally, it is worth investigating whether channel diagonalization is still optimal for EE optimization subject to the deterministic CSI error. To the author's best knowledge, this important issue has not been addressed in the existing literature. For the stochastic CSI error, which is well suited for modeling estimation inaccuracy, existing works typically involve outage-type performance optimization [20], [21]. In this context, the system design is generally more difficult than that for the deterministic CSI error. There exist only a few works considering mean EE optimization under outage constraints, in which the channel diagonalization is generally unavailable [21].

This paper mainly investigates the robust EE maximization of MIMO AF relaying systems under imperfect CSI with the deterministically bounded CSI error. We aim to jointly optimizing the source covariance matrix and relay beamforming matrix/matrices to maximize the system's EE for MIMO AF relaying networks under both additive and multiplicative CSI errors. The main challenge of this robust EE maximization design is that it is essentially a two-objective optimization, namely, maximizing the achievable worst-case rate while minimizing the total power consumption. Clearly, these two objectives are conflicting. Therefore, the optimal diagonalization transmission strategy for traditional robust capacity maximization [17] is not applicable, since both the achievable worst-case rate and the total transmit power are simultaneously maximized by diagonalizing the channel matrices. The robust EE maximization design must find an optimal trade off between maximizing the worst-case rate and minimizing the total power consumption. Our main contributions are summarized as follows:

- Under the spectral-norm constrained additive and multiplicative relay-destination channel errors, we prove the existence of a saddle point for the robust max-min EE problem. The analytical structure of this saddle point is derived, which enables the scalarized reformulation of the robust EE problem to reduce the optimization complexity remarkably. An important insight provided by the optimal analytical solutions is that the eigenmode transmission strategy is optimal for the robust EE maximization under the deterministic CSI errors. We further show that this eigenmode transmission strategy remains optimal for the robust EE optimization when the statistically imperfect source-relay channel is additionally imposed.
- In order to effectively solve the scalarized EE optimization, we propose an alternating optimization of the source covariance matrix related subproblem and the relay beamforming related subproblem. For both these subproblems, we can jointly apply Dinkelbach's method [22] and the Lagrangian dual method to obtain the water-filling structured solutions. The convergence of the proposed alternating optimization is established. This approach is also applicable to the case with additional statistically imperfect source-relay channel.

• Furthermore, we extend our work to multi-hop MIMO AF relaying scenarios. The eigenmode transmission strategy is also proved to be optimal for the robust EE optimization subject to deterministic relay-destination CSI errors, and our proposed alternating optimization remains applicable. The extension to the robust EE optimization subject to additional statistically imperfect relay-relay channels is also discussed.

The bold-faced lower-case and upper-case letters stand for vectors and matrices, respectively. The transpose, Hermitian and inverse operators are denoted by $(\cdot)^{T}$, $(\cdot)^{H}$ and $(\cdot)^{-1}$, respectively, while Tr(A) and det(A) denote the trace and determinant of A, respectively. $E[\cdot]$ is the expectation, and I_n is the $n \times n$ identity matrix, while $\|\cdot\|_2$ denotes the matrix spectral norm, and $A \succeq 0$ indicates that the square matrix A is positive semidefinite. $\mathbf{0}_{n \times m}$ and $\mathbf{1}_n$ denote the $n \times m$ zero matrix and the *n*-dimensional vector with all elements being one, respectively. The $n \times n$ square diagonal matrix with the diagonal elements a_1, a_2, \cdots, a_n is denoted by diag $\{a_1, a_2, \dots, a_n\}$, and similarly for the $m \times n$ diagonal rectangular matrix, all the off-diagonal elements are zero. $A \otimes B$ represents either A or B depending on which one is actually considered. The rank of A is denoted by rank(A), and $(a)^+ = \max\{a, 0\}$. The words 'independently and identically distributed' and 'with respect to' are abbreviated as 'i.i.d.' and 'w.r.t.', respectively.

II. SYSTEM MODEL AND PROBLEM FORMULATION

A. Two-Hop MIMO AF Relaying Networks

Consider a MIMO AF relaying network consisting of an N_S -antenna source, an N_R -antenna relay and an N_D -antenna destination, which operates in half-duplex mode. In the first hop, the source transmits the data vector $s \in \mathbb{C}^{N_S}$ having the covariance matrix $\mathsf{E}[ss^{\mathrm{H}}] = V_S \in \mathbb{C}^{N_S \times N_S}$ to the relay, whose received signal $y_R \in \mathbb{C}^{N_R}$ is expressed as

$$\boldsymbol{y}_R = \boldsymbol{H}_{SR}\boldsymbol{s} + \boldsymbol{n}_R, \tag{1}$$

where $n_R \in \mathbb{C}^{N_R}$ is the additive white Gaussian noise (AWGN) vector of the source-relay link with the covariance matrix $\sigma_r^2 I_{N_R}$, and $H_{SR} \in \mathbb{C}^{N_R \times N_S}$ is the source-relay channel matrix. The source's transmit signal *s* has the power $P_S = \text{Tr}(V_S)$. In the second hop, the relay retransmits the signal received in the first hop by pre-multiplying y_R with the AF beamforming matrix $W_R \in \mathbb{C}^{N_R \times N_R}$. Thus, the relay's transmitted signal $y'_R = W_R y_R$ has the power

$$P_{R} = \operatorname{Tr}\left(\boldsymbol{W}_{R}\left(\boldsymbol{H}_{SR}\boldsymbol{V}_{S}\boldsymbol{H}_{SR}^{\mathrm{H}} + \sigma_{r}^{2}\boldsymbol{I}_{N_{R}}\right)\boldsymbol{W}_{R}^{\mathrm{H}}\right).$$
(2)

The signal $\boldsymbol{y}_D \in \mathbb{C}^{N_D}$ received at the destination is then given by

$$\boldsymbol{y}_D = \boldsymbol{H}_{RD} \boldsymbol{W}_R \boldsymbol{H}_{SR} \boldsymbol{s} + \boldsymbol{H}_{RD} \boldsymbol{W}_R \boldsymbol{n}_R + \boldsymbol{n}_D, \quad (3)$$

where $\boldsymbol{H}_{RD} \in \mathbb{C}^{N_D \times N_R}$ is the relay-destination channel matrix, and $\boldsymbol{n}_D \in \mathbb{C}^{N_D}$ is the AWGN vector of the relay-destination channel with the covariance matrix $\sigma_d^2 \boldsymbol{I}_{N_D}$.

We adopt the EE metric of the MIMO AF relaying network as the optimization objective, which is defined as the ratio of the maximum achievable data rate to the total power consumed. The maximum achievable rate or capacity measured in [bit/s] is expressed as

$$R_{\rm D} = \frac{B}{2} \text{logdet} \left(\boldsymbol{I}_{N_D} + \boldsymbol{H}_{RD} \boldsymbol{W}_R \boldsymbol{H}_{SR} \boldsymbol{V}_S \boldsymbol{H}_{SR}^{\rm H} \boldsymbol{W}_R^{\rm H} \boldsymbol{H}_{RD}^{\rm H} \right)$$
$$\times \cdot \left(\sigma_r^2 \boldsymbol{H}_{RD} \boldsymbol{W}_R \boldsymbol{W}_R^{\rm H} \boldsymbol{H}_{RD}^{\rm H} + \sigma_d^2 \boldsymbol{I}_{N_D} \right)^{-1} \right), \quad (4)$$

where B denotes the allocated system bandwidth and the factor $\frac{1}{2}$ indicates the half-duplex loss. We model the total power consumption of the relaying network as the sum of the transmission powers of the source and relay, scaled by their respective power amplifier efficiencies, and the total circuit power consumption P_C , given by

$$\mathcal{P}(\boldsymbol{V}_{S}, \boldsymbol{W}_{R}) = \frac{P_{S}}{\tau_{s}} + \frac{P_{R}}{\tau_{r}} + P_{C} \text{ [Joule/s]}, \quad (5)$$

where $0 < \tau_s \leq 1$ and $0 < \tau_r \leq 1$ are the source and relay power amplifier efficiencies, respectively. According to [10], [23]–[25], the total circuit power consumption P_C can be modeled as $P_C = N_S P_{dy,s} + N_R P_{dy,r} + P_{st}$, where $P_{dy,s}$ and $P_{dy,r}$ are the dynamic power consumption of each RF chain of the source and relay, respectively, while $P_{st} = P_{st,s} + P_{st,r}$ is the total static power overhead of the source and relay, including baseband processing, power supply and cooling power consumption. Reception generally consumes less circuit power than transmission [23]. Therefore, we neglect the circuit power consumption at the destination. From (4) and (5), the EE metric is defined as

$$\operatorname{EE}(\boldsymbol{V}_{S}, \boldsymbol{W}_{R}) = \frac{R_{\mathrm{D}}}{\mathcal{P}(\boldsymbol{V}_{S}, \mathbf{W}_{R})} \text{ [bit/Joule]}. \tag{6}$$

B. Robust EE Optimization Problem

At high signal-to-noise ratio conditions and with optimal pilot design, receiver can acquire an accurate CSI with training [26]. It is reasonable to assume that the CSI is perfectly available at receiver [11], [16], [17]. But transmitter can only acquire this estimated CSI through a finite-rate feedback channel, which introduces the quantization and feedback delay errors [16]. Consequently, the CSI at transmitter is inherently imperfect. Similar to most of the existing literature [15]–[17], we first assume that the relay has the perfect knowledge of the source-relay channel H_{SR} but it can only acquire an imperfect relay-destination channel H_{RD} . However, we also consider the more generic scenario where the perfect knowledge of the source-relay channel is also unavailable. As aforementioned, there exist two different types of imperfect CSI models, the statistically and deterministically imperfect CSI. Different from the works [11], [13], which study the robust EE optimization under the statistically imperfect CSI, we study the robust EE maximization for deterministically imperfect CSI. In general, the deterministic CSI errors take two different forms, additive CSI errors and multiplicative CSI errors. According to [18], [19], the CSI feedback and quantization errors are considered to be additive, while CSI calibration mismatch and the channel dynamic variations are regarded as

multiplicative errors. By applying the two types of CSI errors to the relay-destination channel H_{RD} , we have

$$\boldsymbol{H}_{RD} = \begin{cases} \widehat{\boldsymbol{H}}_{RD} + \boldsymbol{\Delta}_{RD}, & \|\boldsymbol{\Delta}_{RD}\|_{2} \leq \epsilon_{a}, \\ (\boldsymbol{I}_{N_{D}} + \boldsymbol{E}_{RD}) \widehat{\boldsymbol{H}}_{RD}, & \|\boldsymbol{E}_{RD}\|_{2} \leq \epsilon_{m}, \end{cases}$$
(7)

where $\widehat{\boldsymbol{H}}_{RD} \in \mathbb{C}^{N_D \times N_R}$ is the known nominal relay-destination channel, $\boldsymbol{\Delta}_{RD} \in \mathbb{C}^{N_D \times N_R}$ and $\boldsymbol{E}_{RD} \in \mathbb{C}^{N_D \times N_D}$ are the additive and multiplicative CSI errors, respectively, while ϵ_a and ϵ_m are the corresponding spectral norm bounds of the CSI errors. To focus on the underlying principles and without loss of generality, we consider additive CSI errors and multiplicative CSI errors separately¹. Note that the spectral norm belongs to the unitarily-invariant norm sets, in which the norm-bounded terms are statistically independent and identical in all directions [15]. It also acts as the lower bound of all unitarily-invariant norms. Hence, for the same CSI errors, the spectral norm constrained case covers the largest uncertainty region [15]. Furthermore, when considering another popular Frobenius norm expression, we have $\|\boldsymbol{\Delta}_{RD}\|_2 \leq \|\boldsymbol{\Delta}_{RD}\|_F \leq \sqrt{N_D} \|\boldsymbol{\Delta}_{RD}\|_2$, which indicates that the spectral norm constrained CSI errors can also provide valuable insights for the Frobenius norm constrained case. Given the imperfect CSI specified by (7), the EE metric (6) also depends on $\Delta_{RD} \varnothing E_{RD}$ and, therefore, it is expressed as $\operatorname{EE}(\boldsymbol{V}_S, \boldsymbol{W}_R, \boldsymbol{\Delta}_{RD} \otimes \boldsymbol{E}_{RD}).$

Following the worst case robustness logic, the source covariance matrix V_S and relay beamforming matrix W_R are jointly designed by guaranteeing the maximum EE for all possible relay-destination channel realizations within the uncertainty region defined by (7). This robust EE optimization problem of MIMO AF relaying networks is formulated as

$$\max_{\boldsymbol{V}_{S}, \boldsymbol{W}_{R}} \min_{\boldsymbol{\Delta}_{RD} \otimes \boldsymbol{E}_{RD}} \operatorname{EE} \left(\boldsymbol{V}_{S}, \boldsymbol{W}_{R}, \boldsymbol{\Delta}_{RD} \otimes \boldsymbol{E}_{RD} \right),$$
s.t. $\operatorname{Tr}(\boldsymbol{V}_{S}) \leq P_{S_{\max}},$
 $\operatorname{Tr} \left(\boldsymbol{W}_{R} \left(\boldsymbol{H}_{SR} \boldsymbol{V}_{S} \boldsymbol{H}_{SR}^{\mathrm{H}} + \sigma_{r}^{2} \boldsymbol{I}_{N_{R}} \right) \boldsymbol{W}_{R}^{\mathrm{H}} \right) \leq P_{R_{\max}},$
 $\| \boldsymbol{\Delta}_{RD} \|_{2} \leq \epsilon_{a} \text{ or } \| \boldsymbol{E}_{RD} \|_{2} \leq \epsilon_{m},$
(8)

where $P_{S_{\text{max}}}$ and $P_{R_{\text{max}}}$ are the maximum transmit powers of source and relay, respectively. As (8) contains the interrelated optimization variables and the semi-infinite CSI errors, the classical saddle point theory for concave-convex problems [27, Theorem 36.3] cannot be applied. According to [9], the max-min EE problem (8) can be reduced to a NP-hard sigmoidal programming [28, Theorem 1, page 15]. Therefore, it is also NP-hard and very difficult to solve directly.

III. WORST CASE EE MAXIMIZATION FOR TWO-HOP MIMO AF RELAYING

A. Derivation of Saddle Point

To simplify the intricate relationships among the optimization variables $\{V_S, W_R, \Delta_{RD} \otimes E_{RD}\}$, we utilize the

¹Our work can easily be extended to the case having both additive and multiplicative CSI errors, namely, $H_{RD} = (I_{N_D} + E_{RD})\hat{H}_{RD} + \Delta_{RD}$. This is because in this case, similar worst-case channel-diagonalizing structure can easily be derived by applying the results of this work for the additive CSI errors (Δ_{RD}) and the multiplicative CSI errors ($E_{RD}\hat{H}_{RD}$).

Woodbury matrix identity to equivalently transform (4) into

$$R_{\rm D} = \frac{B}{2} \log \det \left(\boldsymbol{I}_{N_R} + \sigma_r^{-2} \boldsymbol{H}_{SR} \boldsymbol{V}_S \boldsymbol{H}_{SR}^{\rm H} - \sigma_r^{-2} \boldsymbol{H}_{SR} \boldsymbol{V}_S \boldsymbol{H}_{SR}^{\rm H} \right. \\ \times \left(\boldsymbol{I}_{N_R} + \sigma_r^2 \sigma_d^{-2} \boldsymbol{W}_R^{\rm H} \boldsymbol{H}_{RD}^{\rm H} \boldsymbol{H}_{RD} \boldsymbol{W}_R \right)^{-1} \right) \\ = \mathcal{R} \left(\boldsymbol{V}_S, \boldsymbol{W}_R, \boldsymbol{\Delta}_{RD} \otimes \boldsymbol{E}_{RD} \right).$$
(9)

Then the EE metric in (8) is rewritten as EE $(V_S, W_R, \Delta_{RD} \otimes E_{RD}) = \mathcal{R}(V_S, W_R, \Delta_{RD} \otimes E_{RD}) / \mathcal{P}(V_S, W_R)$. Since (8) is not concave-convex in $\{V_S, W_R, \Delta_{RD} \otimes E_{RD}\}$, it is difficult to solve it directly. Instead, we consider its counterpart, i.e., the following min-max EE problem,

$$\begin{array}{l} \min_{\boldsymbol{\Delta}_{RD} \otimes \boldsymbol{E}_{RD}} \quad \max_{\boldsymbol{V}_{S}, \boldsymbol{W}_{R}} \operatorname{EE}\left(\boldsymbol{V}_{S}, \boldsymbol{W}_{R}, \boldsymbol{\Delta}_{RD} \otimes \boldsymbol{E}_{RD}\right) \\ \text{s.t.} \quad \operatorname{Tr}(\boldsymbol{V}_{S}) \leq P_{S_{\max}}, \\ \quad \operatorname{Tr}\left(\boldsymbol{W}_{R}\left(\boldsymbol{H}_{SR} \boldsymbol{V}_{S} \boldsymbol{H}_{SR}^{\mathrm{H}} + \sigma_{r}^{2} \boldsymbol{I}_{N_{R}}\right) \boldsymbol{W}_{R}^{\mathrm{H}}\right) \leq P_{R_{\max}}, \\ \quad \|\boldsymbol{\Delta}_{RD}\|_{2} \leq \epsilon_{a} \text{ or } \|\boldsymbol{E}_{RD}\|_{2} \leq \epsilon_{m}. \quad (10)
\end{array}$$

Generally, $\max_{x} \min_{y} f(x,y) \leq \min_{y} \max_{x} f(x,y)$ [29, Section 5.4] holds implying that the max-min problem and the min-max problem are not identical. However, according to [30, Corollary 9.16], if there exists a saddle point for EE($V_S, W_R, \Delta_{RD} \otimes E_{RD}$), then it is globally optimal for both problems (8) and (10). Hence, we first study the problem (10) and derive its optimal solution, and then prove that the obtained solution is indeed a saddle point of EE($V_S, W_R, \Delta_{RD} \otimes E_{RD}$). Let's define the singular value decomposition (SVD) of H_{SR} and \widehat{H}_{RD} as

$$\boldsymbol{H}_{SR} = \boldsymbol{U}_{SR} \boldsymbol{\Sigma}_{SR} \boldsymbol{Q}_{SR}^{\mathrm{H}}, \qquad (11)$$

$$\widehat{\boldsymbol{H}}_{RD} = \widehat{\boldsymbol{U}}_{RD} \widehat{\boldsymbol{\Sigma}}_{RD} \widehat{\boldsymbol{Q}}_{RD}^{"}, \qquad (12)$$

where $U_{SR} \in \mathbb{C}^{N_R \times N_R}$ and $Q_{SR} \in \mathbb{C}^{N_S \times N_S}$ as well as $\widehat{U}_{RD} \in \mathbb{C}^{N_D \times N_D}$ and $\widehat{Q}_{RD} \in \mathbb{C}^{N_R \times N_R}$ are the unitary singular matrices for H_{SR} and \widehat{H}_{RD} , while the diagonal rectangular matrices $\Sigma_{SR} \in \mathbb{C}^{N_R \times N_S}$ and $\widehat{\Sigma}_{RD} \in \mathbb{C}^{N_D \times N_R}$ take $N_P = \min\{N_R, N_S\}$ singular values (SVs) $\{\sigma_{sr,1}, \cdots, \sigma_{sr,N_P}\}$ of H_{SR} and $N_C = \min\{N_D, N_R\}$ SVs $\{\widehat{\sigma}_{rd,1}, \cdots, \widehat{\sigma}_{rd,N_C}\}$ of \widehat{H}_{RD} as diagonal elements.

Theorem 1: For the min-max EE problem (10), the optimal source covariance matrix V_S^* , the optimal relay beamforming matrix W_R^* and the worst-case CSI errors $\Delta_{RD}^* \otimes E_{RD}^*$ satisfy

$$\boldsymbol{V}_{S}^{\star} = \boldsymbol{Q}_{SR} \boldsymbol{\Sigma}_{S} \boldsymbol{Q}_{SR}^{\mathrm{H}}, \tag{13}$$

$$\boldsymbol{W}_{R}^{\star} = \widehat{\boldsymbol{Q}}_{RD} \boldsymbol{\Sigma}_{X} \left(\boldsymbol{I}_{N_{R}} + \sigma_{r}^{-2} \boldsymbol{\Sigma}_{SR} \boldsymbol{\Sigma}_{S} \boldsymbol{\Sigma}_{SR}^{\mathrm{H}} \right)^{-\frac{1}{2}} \boldsymbol{U}_{SR}^{\mathrm{H}}, \quad (14)$$

$$\Delta_{RD}^{\star} = -\widehat{U}_{RD}\Lambda_{RD}\widehat{Q}_{RD}^{n} \text{ or } E_{RD}^{\star} = -\epsilon_{m}I_{N_{D}}, \qquad (15)$$

where $\Sigma_S = \operatorname{diag}\{\lambda_{s,1}, \dots, \lambda_{s,N_S}\}\$ and $\Sigma_X = \operatorname{diag}\{\sigma_{x,1}, \dots, \sigma_{x,N_C}, 0, \dots, 0\} \in \mathbb{C}^{N_R \times N_R}$, in which $\lambda_{s,i}$ for $1 \le i \le N_S$ and $\sigma_{x,j}$ for $1 \le j \le N_C$ replace V_S and W_R as the optimization scalar variables for the min-max EE problem (10), while the diagonal rectangular matrix $\Lambda_{RD} \in \mathbb{C}^{N_D \times N_R}$ has the N_C diagonal elements $\min\{\widehat{\sigma}_{rd,1}, \epsilon_a\}, \dots, \min\{\widehat{\sigma}_{rd,N_C}, \epsilon_a\}$.

Proof: See Appendix A. \Box

Theorem 2: The optimal solution $\{V_S^*, W_R^*, \Delta_{RD}^* \otimes E_{RD}^*\}$ of the min-max EE problem (10) provided by Theorem 1 is the

saddle point of the EE metric $\text{EE}(V_S, W_R, \Delta_{RD} \otimes E_{RD})$, *i.e.*,

$$EE(\boldsymbol{V}_{S}, \boldsymbol{W}_{R}, \boldsymbol{\Delta}_{RD}^{\star} \otimes \boldsymbol{E}_{RD}^{\star}) \leq EE(\boldsymbol{V}_{S}^{\star}, \boldsymbol{W}_{R}^{\star}, \boldsymbol{\Delta}_{RD}^{\star} \otimes \boldsymbol{E}_{RD}^{\star})$$

$$\leq EE(\boldsymbol{V}_{S}^{\star}, \boldsymbol{W}_{R}^{\star}, \boldsymbol{\Delta}_{RD} \otimes \boldsymbol{E}_{RD}), \quad (16)$$

holds for any feasible V_S , W_R and $\Delta_{RD} \otimes E_{RD}$. According to [30, Corollary 9.16], it is also optimal for the original max-min EE problem (8).

Proof: See Appendix B.
$$\Box$$

According to Theorems 1 and 2, the optimal V_S^{\star} , W_R^{\star} and $\Delta_{RD}^{\star} \otimes E_{RD}^{\star}$ that solve the max-min EE problem (8) all have the channel-diagonalizing structure. Calculating the optimal V_S^{\star} and W_R^{\star} becomes determining the values of $\lambda_s = [\lambda_{s,1} \cdots \lambda_{s,N_S}]^{\mathrm{T}}$ and $\sigma_x = [\sigma_{x,1}^2 \cdots \sigma_{x,N_C}^2]^{\mathrm{T}}$.

B. Proposed Alternating Optimization Algorithm

Based on Theorem 1, the original max-min EE problem (8) with matrix variables can be equivalently transformed into the problem (17) with scalar variables, shown at the top of the next page, where $N_L = \min \{N_P, N_C\}$ and for $1 \le i \le N_C$,

$$\widetilde{\sigma}_{rd,i} = \begin{cases} \left(\widehat{\sigma}_{rd,i} - \epsilon_a\right)^+ & \text{for additive CSI errors,} \\ (1 - \epsilon_m)^+ \widehat{\sigma}_{rd,i} & \text{for multiplicative CSI errors,} \end{cases}$$
(18)

Compared to the original problem (8), the number of optimization variables in the problem (17) is significantly reduced, namely, from $N_S^2 + N_R^2 + N_D N_R$ to $N_S + N_C$. In order to efficiently solve the problem (17), the fractional programming theory [22], [30] is first introduced.

Lemma 1: ([22], [30]) Given a fractional function $f(\mathbf{A}) = \frac{N(\mathbf{A})}{G(\mathbf{A})}$, provided that $N(\mathbf{A})$ and $G(\mathbf{A})$ are concave and convex w.r.t. \mathbf{A} , respectively, then $f(\mathbf{A})$ is quasi-concave. By introducing an auxiliary variable η , a single-parameter subtractive function is defined as

$$F(\eta) = \max_{\boldsymbol{A}} N(\boldsymbol{A}) - \eta G(\boldsymbol{A}).$$
(19)

The inner maximization problem of (19) is concave w.r.t. A for any fixed η , and $F(\eta)$ is a decreasing function of η . Moreover, the problem of maximizing $f(\mathbf{A})$ is equivalent to finding the zero point of $F(\eta)$, and Dinkelbach's method can be invoked for finding $F(\eta) = 0$, which is guaranteed to converge to a globally optimal solution of maximizing $f(\mathbf{A})$ [22].

Taking the second derivative of the numerator of the objective function in (17) w.r.t. λ_s for fixed σ_x , it is seen that the numerator of the objective function is concave w.r.t. λ_s . The denominator of the objective function in (17) is linear w.r.t. λ_s given σ_x . According to Lemma 1, the problem (17) is quasi-concave for λ_s given σ_x . Similarly, the problem (17) is quasi-concave for σ_x given λ_s . Therefore, it can be efficiently tackled by an alternating optimization between the subproblem of optimizing λ_s for fixed σ_x and that of optimizing σ_x for fixed λ_s .

By introducing the auxiliary variable η based on Lemma 1, we transform (17) into the following single-parameter

$$\max_{\boldsymbol{\lambda}_{s},\boldsymbol{\sigma}_{x}} \frac{\sum_{i=1}^{N_{L}} \log \left(\frac{1 + \sigma_{r}^{2} \sigma_{d}^{2} \sigma_{x,i}^{2} \tilde{\sigma}_{d,i}^{2}}{1 + \sigma_{r}^{-2} \sigma_{sr,i}^{2} \lambda_{s,i} + \sigma_{r}^{2} \sigma_{d}^{-2} \sigma_{x,i}^{2} \tilde{\sigma}_{rd,i}^{2}} \right) + \sum_{i=1}^{N_{L}} \log \left(1 + \sigma_{r}^{-2} \sigma_{sr,i}^{2} \lambda_{s,i} \right)}{\sum_{i=1}^{N_{S}} \lambda_{s,i} + \sum_{i=1}^{N_{C}} \sigma_{r}^{2} \sigma_{x,i}^{2} + P_{C}},$$
s.t.
$$\sum_{i=1}^{N_{S}} \lambda_{s,i} \leq P_{S_{max}}, \quad \sum_{i=1}^{N_{C}} \sigma_{r}^{2} \sigma_{x,i}^{2} \leq P_{R_{max}}, \quad (17)$$

subtractive problem

$$\max_{\boldsymbol{\lambda}_{s},\boldsymbol{\sigma}_{x}} \sum_{i=1}^{N_{L}} \left(\log \left(\frac{1 + \sigma_{r}^{2} \sigma_{d}^{-2} \sigma_{x,i}^{2} \widetilde{\sigma}_{rd,i}^{2}}{1 + \sigma_{r}^{-2} \sigma_{sr,i}^{2} \lambda_{s,i} + \sigma_{r}^{2} \sigma_{d}^{-2} \sigma_{x,i}^{2} \widetilde{\sigma}_{rd,i}^{2}} \right) \\
+ \log \left(1 + \sigma_{r}^{-2} \sigma_{sr,i}^{2} \lambda_{s,i} \right) \right) \\
- \eta \left(\sum_{i=1}^{N_{S}} \lambda_{s,i} + \sum_{i=1}^{N_{C}} \sigma_{r}^{2} \sigma_{x,i}^{2} + P_{C} \right), \\
\text{s.t.} \quad \sum_{i=1}^{N_{S}} \lambda_{s,i} \leq P_{S_{\max}}, \sum_{i=1}^{N_{C}} \sigma_{r}^{2} \sigma_{x,i}^{2} \leq P_{R_{\max}}. \quad (20)$$

For given η , since (20) is strictly concave w.r.t λ_s for fixed σ_x and vice versa, the Lagrangian dual method can be adopted for obtaining the corresponding optimal solutions to the respective subproblems. Specifically, given η , we introduce the Lagrangian dual function of (20) as

$$L(\boldsymbol{\lambda}_{s}, \boldsymbol{\sigma}_{x}, \boldsymbol{\mu}, \boldsymbol{\beta}; \boldsymbol{\eta}) = \sum_{i=1}^{N_{L}} \log \left(\frac{1 + \sigma_{r}^{2} \sigma_{d}^{-2} \sigma_{x,i}^{2} \widetilde{\sigma}_{rd,i}^{2}}{1 + \sigma_{r}^{-2} \sigma_{sr,i}^{2} \lambda_{s,i} + \sigma_{r}^{2} \sigma_{d}^{-2} \sigma_{x,i}^{2} \widetilde{\sigma}_{rd,i}^{2}} \right) + \sum_{i=1}^{N_{L}} \log \left(1 + \sigma_{r}^{-2} \sigma_{sr,i}^{2} \lambda_{s,i} \right) - (\boldsymbol{\eta} + \boldsymbol{\mu}) \sum_{i=1}^{N_{S}} \lambda_{s,i} - (\boldsymbol{\eta} + \boldsymbol{\beta}) \sum_{i=1}^{N_{C}} \sigma_{r}^{2} \sigma_{x,i}^{2} + C,$$
(21)

where μ and β are the Lagrangian multipliers for the source and relay transmit power constraints, respectively, and $C = \mu P_{S_{\text{max}}} + \beta P_{R_{\text{max}}} - \eta P_C$. For the subproblem of optimizing σ_x given λ_s , we take the derivative of $L(\lambda_s, \sigma_x, \mu, \beta; \eta)$ w.r.t. $\sigma_{x,i}$ and utilize Karush-Kuhn-Tucker (KKT) conditions to obtain the optimal σ_x for fixed λ_s , denoted by $\sigma_x(\lambda_s; \eta)$, as (22), shown at the bottom of the next page. where β satisfies $\beta \left(\sum_{i=1}^{N_C} \sigma_r^2 \sigma_{x,i}^2 - P_{R_{\text{max}}} \right) = 0$ and it can be determined by the bisection search owing to the monotonically decreasing property of $\sigma_{x,i}^2(\lambda_s; \eta)$ w.r.t. β . Similarly, for fixed σ_x , we have the optimal λ_s , denoted by $\lambda_s(\sigma_x; \eta)$, in (23), shown at the bottom of the next page. where owing to the monotonically decreasing property of $\lambda_s(\sigma_x; \eta)$ w.r.t μ , μ is chosen by the bisection search to ensure $\mu \left(\sum_{i=1}^{N_S} \lambda_{s,i} - P_{S_{\text{max}}} \right) = 0$.

For efficiently realizing the worst-case EE maximization, we apply Dinkelbach's method to both the subproblems of (20) to update η by utilizing the optimal σ_x in (22) for fixed λ_s and by utilizing the optimal λ_s in (23) for fixed σ_x , respectively. Dinkelbach's method is an iterative optimization process, which converges when the zero objective value of the problem (20) is realized. Specifically, the update of η in Dinkelbach's method based on $\{\lambda_s, \sigma_x\}$ is given by (24), shown at the bottom of the next page. Integrating (22) to (24), the proposed alternating optimization for the worst-case EE maximization under additive or multiplicative CSI errors is summarized in Algorithm 1.

C. Convergence and the Speedup Strategy

For characterizing the convergence, let us consider the arbitrary feasible initial value $\lambda_s^{(n)}$ and $\eta^{(n)}$ for the nth iteration of Algorithm 1. According to Lemma 1, given $\lambda_s^{(n)}$, the σ_x related subproblem of (20) is strictly concave, and we can obtain the unique and globally optimal $\sigma_x^{(n+1)}$ in (22), which implies that after Step 2 of Algorithm 1, $\operatorname{EE}(\boldsymbol{V}_{S}^{(n)}, \boldsymbol{W}_{R}^{(n)}, \boldsymbol{\Delta}_{RD}^{\star} \otimes \boldsymbol{E}_{RD}^{\star}) \leq$ $\operatorname{EE}(\boldsymbol{V}_{S}^{(n)}, \boldsymbol{W}_{R}^{(n+1)}, \boldsymbol{\Delta}_{RD}^{\star} \otimes \boldsymbol{E}_{RD}^{\star}). \quad \text{Given } \boldsymbol{\sigma}_{x}^{(n+1)}, \text{ the } \boldsymbol{\lambda}_{s}$ related subproblem of (20) is also concave, and the globally optimal and unique $\lambda_s^{(n+1)}$ is derived in (23), which implies that after Step 3 of Algorithm 1, we have $\operatorname{EE}(\boldsymbol{V}_{S}^{(n)}, \boldsymbol{W}_{R}^{(n+1)})$ $\boldsymbol{\Delta}_{RD}^{\star} \otimes \boldsymbol{E}_{RD}^{\star} \right) \leq \operatorname{EE} \left(\boldsymbol{V}_{S}^{(n+1)}, \boldsymbol{W}_{R}^{(n+1)}, \boldsymbol{\Delta}_{RD}^{\star} \otimes \boldsymbol{E}_{RD}^{\star} \right). \quad \operatorname{Com-}$ bining these two non-strict inequalities, we generally have $\text{EE}(\boldsymbol{V}_{S}^{(n)}, \boldsymbol{W}_{R}^{(n)}, \boldsymbol{\Delta}_{RD}^{\star} \otimes \boldsymbol{E}_{RD}^{\star}) < \text{EE}(\boldsymbol{V}_{S}^{(n+1)}, \boldsymbol{W}_{R}^{(n+1)}, \boldsymbol{\Delta}_{RD}^{\star} \otimes \boldsymbol{E}_{RD}^{\star})$ after the (n + 1)th iteration. It is widely exploited that the worst-case robust maximization is generally upper bounded by the corresponding perfect-case maximization [11], [13]. By setting $\sigma_{rd,i} = \hat{\sigma}_{rd,i}, 1 \leq i \leq N_C$, in the problem (20) to indicate that the relay's knowledge of H_{RD} is perfect, it becomes the perfect-case EE maximization, which has been solved in [11] and the solution provides an effective upper bound for our worst-case EE. Since the achievable worst-case EE of Algorithm 1 is non-decreasing and upper bounded, we conclude that Algorithm 1 is guaranteed to converge to a stationary point $\{\lambda_s^{\text{st}}, \sigma_x^{\text{st}}\}$ of the problem (17).

Substituting this stationary point of the problem (17) into (13) and (14) of Theorem 1 yields a locally optimal solution $\{V_S^{\text{st}}, W_R^{\text{st}}\}$ to the problem (8) $\begin{array}{l} \text{given} \quad \boldsymbol{\Delta}_{RD}^{\star} \otimes \boldsymbol{E}_{RD}^{\star}, \quad \text{which} \quad \text{satisfies} \quad \mathrm{EE}(\boldsymbol{V}_{S}^{\mathrm{st}}, \boldsymbol{W}_{R}^{\mathrm{st}}, \\ \boldsymbol{\Delta}_{RD}^{\star} \otimes \boldsymbol{E}_{RD}^{\star}) \geq \quad \mathrm{EE}(\boldsymbol{V}_{S}, \boldsymbol{W}_{R}, \boldsymbol{\Delta}_{RD}^{\star} \otimes \boldsymbol{E}_{RD}^{\star}) \quad \text{and} \\ (\boldsymbol{V}_{S}, \boldsymbol{W}_{R}) \in \mathcal{U}(\boldsymbol{V}_{S}^{\mathrm{st}}, \boldsymbol{W}_{R}^{\mathrm{st}}; \delta_{d}) \quad (\text{Here,} \quad \mathcal{U}(\boldsymbol{V}_{S}^{\mathrm{st}}, \boldsymbol{W}_{R}^{\mathrm{st}}; \delta_{d}) \end{aligned}$ denotes the vicinity sphere of point $\{V_S^{\mathrm{st}}, W_R^{\mathrm{st}}\}$ with radius δ_d). Moreover, given $\{V_S^{st}, W_R^{st}\}$, it is observed that the worst-case additive/multiplicative error to the problem (8) is still $\Delta_{RD}^{\star} \otimes E_{RD}^{\star}$, and thus we have
$$\begin{split} & \operatorname{EE}(\boldsymbol{V}_{S}^{\operatorname{st}}, \boldsymbol{W}_{R}^{\operatorname{st}}, \boldsymbol{\Delta}_{RD}^{\star} \otimes \boldsymbol{E}_{RD}^{\star}) \leq \operatorname{EE}(\boldsymbol{V}_{S}^{\operatorname{st}}, \boldsymbol{W}_{R}^{\operatorname{st}}, \boldsymbol{\Delta}_{RD} \otimes \boldsymbol{E}_{RD}). \\ & \text{Therefore, for any } (\boldsymbol{V}_{S}, \boldsymbol{W}_{R}) \in \mathcal{U}(\boldsymbol{V}_{S}^{\operatorname{st}}, \boldsymbol{W}_{R}^{\operatorname{st}}; \delta_{d}), \end{split}$$
 $\text{it yields } \operatorname{EE} \bigl(\boldsymbol{V}_S, \boldsymbol{W}_R, \hspace{0.1 in} \boldsymbol{\Delta}^{\star}_{RD} \varnothing \boldsymbol{E}^{\star}_{RD} \bigr) \hspace{0.1 in} \leq \hspace{0.1 in} \operatorname{EE} \bigl(\boldsymbol{V}^{\mathrm{st}}_S, \hspace{0.1 in}$ $oldsymbol{W}_R^{ ext{st}}, oldsymbol{\Delta}_{RD}^{\star} arnothing oldsymbol{E}_{RD} ig) \ \leq \ ext{EE}ig(oldsymbol{V}_S^{ ext{st}}, oldsymbol{W}_R^{ ext{st}}, oldsymbol{\Delta}_{RD}^{ ext{st}} arnothing oldsymbol{E}_{RD}ig), \ ext{based}$ based on which we can conclude that $\{V_S^{\text{st}}, W_R^{\text{st}} \mid \Delta_{RD}^{\star} \otimes E_{RD}^{\star}\}$ is a local saddle point (stationary point) of the maxmin EE problem (8). In addition, it is readily inferred that by setting $\eta = \eta^{(n)}$ and $\eta = \tilde{\eta}$ for Step 2 and Step 3, respectively, the objective function value of

Algorithm 1 The proposed alternating optimization for solving (17)

Require: The initial $\lambda_s^{(0)}$ and $\eta^{(0)}$; a sufficiently small tolerance threshold $\zeta > 0$; the iteration index n = 0;

1: repeat

- 2: Fix $\lambda_s = \lambda_s^{(n)}$, start from $\eta = \eta^{(n)}$, apply Dinkelbach method to iteratively optimize between $\sigma_x(\lambda_s^{(n)};\eta)$ of (22) and $\eta(\lambda_s^{(n)},\sigma_x)$ of (24) to obtain $\sigma_x^{(n+1)}$ and $\tilde{\eta}$ that realize the zero objective value of (20);
- 3: Fix $\sigma_x = \sigma_x^{(n+1)}$, start from $\eta = \tilde{\eta}$, apply Dinkelbach method to iteratively optimize between $\lambda_s(\sigma_x^{(n+1)};\eta)$ of (23) and $\eta(\lambda_s, \sigma_x^{(n+1)})$ of (24) to obtain $\lambda_s^{(n+1)}$ and $\eta^{(n+1)}$ that realize the zero objective value of (20);

4:
$$n = n + 1;$$

5: **until** $|\operatorname{EE}(\boldsymbol{V}_{S}^{(n)}, \boldsymbol{W}_{R}^{(n)}, \boldsymbol{\Delta}_{RD}^{(n)} \otimes \boldsymbol{E}_{RD}^{(n)})| \leq \zeta;$ $\operatorname{EE}(\boldsymbol{V}_{S}^{(n-1)}, \boldsymbol{W}_{R}^{(n-1)}, \boldsymbol{\Delta}_{RD}^{(n-1)} \otimes \boldsymbol{E}_{RD}^{(n-1)})| < \zeta;$

Ensure:
$$\operatorname{EE}(V_{S}^{\star}, W_{R}^{\star}, \Delta_{RD}^{\star} \otimes E_{RD}^{\star})$$
 with $V_{S}^{\star} = V_{S}^{(n)}$
 $W_{R}^{\star} = W_{R}^{(n)}$ and $\Delta_{RD}^{\star} \otimes E_{RD}^{\star} = \Delta_{RD}^{(n)} \otimes E_{RD}^{(n)}$;

problem (17) is also non-decreasing between the inner iteration of Step 2 and that of Step 3. In other words, with the obtained $\{W_R^{(n)}, V_S^{(n)}\}$ after the *n*th outer iteration, we further have $\text{EE}(V_S^{(n)}, W_R^{(n)}, \Delta_{RD}^* \otimes E_{RD}^*) \leq \cdots \leq \text{EE}(V_S^{(n)}, W_R^{(n_{in,k})}, \Delta_{RD}^* \otimes E_{RD}^*) \leq \cdots \leq \text{EE}(V_S^{(n)}, W_R^{(n+1)}, \Delta_{RD}^* \otimes E_{RD}^*) \leq \cdots \leq \text{EE}(V_S^{(n+1)}, \Delta_{RD}^* \otimes E_{RD}^*) \leq \cdots \leq \text{EE}(V_S^{(n+1)}, \Delta_{RD}^* \otimes E_{RD}^*) \leq \cdots \leq \text{EE}(V_S^{(n+1)}, M_R^{(n+1)}, \Delta_{RD}^* \otimes E_{RD}^*) \leq \cdots \leq \text{EE}(V_S^{(n+1)}, W_R^{(n+1)}, \Delta_{RD}^* \otimes E_{RD}^*) \leq \cdots \leq \text{EE}(V_S^{(n+1)}, W_R^{(n+1)}, \Delta_{RD}^* \otimes E_{RD}^*)$, where $W_R^{(n_{in,k})}$ and $V_S^{(n_{in,k})}$ denote the optimized W_R and V_S after the *k*th inner iteration of Step 2 and Step 3 in the *n*+1th outer iteration, respectively. Overall, by defining $\eta = \eta^{(n)}$ and $\eta = \tilde{\eta}$ for Step 2 and Step 3 of Algorithm 1, respectively, the achieved EE value for the problem (17) is guaranteed to be non-decreasing both in the inner iteration of Step 2/Step 3 and in the outer iteration between Step 2 and Step 3, which is beneficial for speeding up the convergence of Algorithm 1.

D. Optimality and Complexity

Although a local stationary point $\{\lambda_s^{st}, \sigma_x^{st}\}$ of the problem (17) found by Algorithm 1 is not necessary an optimal solution, it can be seen from the above discussions that the solution $\{V_S^{\mathrm{st}}, W_R^{\mathrm{st}}, \Delta_{RD}^{\star} \otimes E_{RD}^{\star}\}$ associated with $\{\lambda_s^{\rm st}, \sigma_r^{\rm st}\}$ as in Theorem 1 is also a local saddle point to the max-min EE problem (8). A further advantage of Algorithm 1 for solving the non-convex problem (17) is its low computational complexity, on the order of $\mathcal{O}(I_{\text{ite}}(N_S \log_2(N_S)))$ $+N_C \log_2(N_C))$ due to water-filling solution in each step, where I_{ite} denotes the total number of outer-inner iterations. As presented in Section V, Algorithm 1 converges within 10 iterations for both outer and inner loops. This should be contrasted with the computational complexity of brute-force search for finding an optimal solution to the problem (17), namely, $\mathcal{O}\left(\left(\frac{P_{S_{\max}}}{\Delta s}\right)^{N_S}\left(\frac{P_{R_{\max}}}{\Delta s}\right)^{N_C}\right)$, where Δs is the step length. To obtain an accurate solution, a small step length Δs is required, which imposes extremely high complexity. In addition, an upper bound for the objective function of problem (17), denoted as $f_{\rm up}(\boldsymbol{\lambda}_s, \boldsymbol{\sigma}_x)$, is given by (25), shown at the bottom of the next page. The inequality in (25) is derived according to the identity $a+b \ge 2\sqrt{ab}$. Note that this upper bound is jointly quasi-concave w.r.t $\{\lambda_s, \sigma_x\}$, to which the globally optimal solution is available. Moreover, we readily find that this upper bound becomes tight when the sufficiently high transmit powers at source and relay are considered.

To demonstrate the effectiveness of our proposed alternating optimization, Fig. 1 firstly shows the worst-case EE performance as the functions of the source transmit power $P_{S_{\text{max}}}$ for Algorithm 1 and the brute-force search, in comparison with the upper bound (25), where a small-scale system setup with $N_S/N_R/N_D=2/2/2$ and $P_{R_{\text{max}}}=30 \text{ dBm}$ is considered. Then Table I compares the execution times of Algorithm 1 and brute-force search for solving the problem (17). It can

$$\sigma_{x,i}^{2}(\boldsymbol{\lambda}_{s};\eta) = \begin{cases} \frac{\sqrt{\sigma_{r}^{-4}\sigma_{sr,i}^{4}\lambda_{s,i}^{2} + \frac{4\sigma_{d}^{-2}\sigma_{sr,i}^{2}\lambda_{s,i}\tilde{\sigma}_{rd,i}^{2}}{\ln 2(\beta+\eta)\sigma_{r}^{2}} - \sigma_{r}^{-2}\sigma_{sr,i}^{2}\lambda_{s,i} - 2} \\ \frac{2\sigma_{r}^{2}\sigma_{r}^{-2}\tilde{\sigma}_{rd,i}^{2}}{2\sigma_{r}^{2}\sigma_{d}^{-2}\tilde{\sigma}_{rd,i}^{2}}, & \tilde{\sigma}_{rd,i} > 0 \text{ and } 1 \le i \le N_{L}, \\ 0, & \tilde{\sigma}_{rd,i} = 0 \text{ or } N_{L} + 1 \le i \le N_{C}, \end{cases}$$
(22)

$$\lambda_{s,i}(\boldsymbol{\sigma}_{x};\eta) = \begin{cases} \frac{\sqrt{\sigma_{r}^{4}\sigma_{d}^{-4}\sigma_{x,i}^{4}\widetilde{\sigma}_{rd,i}^{4} + \frac{4\sigma_{d}^{-2}\sigma_{sr,i}^{2}\sigma_{x,i}^{2}\widetilde{\sigma}_{rd,i}^{2}}{\ln 2(\mu+\eta)} - \sigma_{r}^{2}\sigma_{d}^{-2}\sigma_{x,i}^{2}\widetilde{\sigma}_{rd,i}^{2} - 2} \\ \frac{2\sigma_{r}^{-2}\sigma_{sr,i}^{2}}{0,} & \sigma_{sr,i} = 0 \text{ or } N_{L} + 1 \le i \le N_{S}, \end{cases}$$
(23)

$$\eta(\boldsymbol{\lambda}_{s},\boldsymbol{\sigma}_{x}) = \frac{\sum_{i=1}^{N_{L}} \log\left(\frac{1 + \sigma_{r}^{2} \sigma_{d}^{2} \sigma_{x,i}^{2} \widetilde{\sigma}_{rd,i}^{2}}{1 + \sigma_{r}^{-2} \sigma_{sr,i}^{2} \lambda_{s,i} + \sigma_{r}^{2} \sigma_{d}^{-2} \sigma_{x,i}^{2} \widetilde{\sigma}_{rd,i}^{2}}\right) + \sum_{i=1}^{N_{L}} \log\left(1 + \sigma_{r}^{-2} \sigma_{sr,i}^{2} \lambda_{s,i}\right)}{\sum_{i=1}^{N_{S}} \lambda_{s,i} + \sum_{i=1}^{N_{C}} \sigma_{r}^{2} \sigma_{x,i}^{2} + P_{C}}.$$
(24)

 TABLE I

 Comparison of Execution Times by the Proposed Alternating Optimization and Brute-Force Search

Time (s) Power(dBm) Method	24	26	28	30	32	34	36	38
Algorithm 1	0.033	0.038	0.061	0.067	0.070	0.071	0.081	0.062
Brute-force search $\Delta s = 0.1$	0.039	0.042	0.0642	0.072	0.147	0.278	1.363	7.558
Brute-force search $\Delta s = 0.05$	0.042	0.057	0.071	0.191	1.076	7.272	110.524	846.584
Brute-force search $\Delta s = 0.01$	1.381	8.116	79.445	744.259	4.874E+3	1.718E+4	5.305E+4	1.599E+5



Fig. 1. The worst-case EE performance as the functions of the source maximum transmit power $P_{S_{\max}}$ for Algorithm 1 and the brute-force search, in comparison with the upper bound.

be seen from Fig. 1 and Table I that there is almost no loss of optimality by using Algorithm 1, which imposes a dramatically lower complexity than the brute-force search. We also observe from Fig. 1 that the gap between the upper bound and optimal solution to the problem (17) is much reduced at high source transmit power.

E. Extension to Imperfect Source-Relay Channel

Since the relay can estimate H_{SR} with higher accuracy, the deterministically imperfect model, such as the one adopted for H_{RD} in (7), is inappropriate for H_{SR} . It is more appropriate to adopt the statistically imperfect CSI model [11], [13], [14] to express H_{SR} as

$$\boldsymbol{H}_{SR} = \boldsymbol{R}_{R}^{\frac{1}{2}} \widetilde{\boldsymbol{H}}_{SR} \boldsymbol{R}_{S}^{\frac{1}{2}}, \qquad (26)$$

where the positive semidefinite $\mathbf{R}_R \in \mathbb{C}^{N_R \times N_R}$ and $\mathbf{R}_S \in \mathbb{C}^{N_S \times N_S}$ are the relay and source spatial correlation matrices, respectively, while $\widetilde{\mathbf{H}}_{SR} \in \mathbb{C}^{N_R \times N_S}$ is a random matrix whose elements are i.i.d. complex Gaussian variables with the distribution $\mathcal{CN}(0, 1)$. Both \mathbf{R}_R and \mathbf{R}_S are available at the relay but the instantaneous $\widetilde{\mathbf{H}}_{SR}$ is unknown. Thus the instantaneous robust EE optimization (8) is infeasible, and the

following average EE metric can be considered [13]

$$\widetilde{\mathsf{EE}}(\boldsymbol{V}_{S}, \boldsymbol{W}_{R}, \boldsymbol{\Delta}_{RD} \otimes \boldsymbol{E}_{RD}) = \frac{\mathsf{E}[R_{\mathrm{D}}(\boldsymbol{V}_{S}, \boldsymbol{W}_{R}, \boldsymbol{\Delta}_{RD} \otimes \boldsymbol{E}_{RD})]}{\mathsf{E}[\mathcal{P}(\boldsymbol{V}_{S}, \boldsymbol{W}_{R})]} = \frac{\mathsf{E}[R_{\mathrm{D}}(\boldsymbol{V}_{S}, \boldsymbol{W}_{R}, \boldsymbol{\Delta}_{RD} \otimes \boldsymbol{E}_{RD})]}{\frac{P_{S}}{T_{T}} + \frac{\mathsf{E}[P_{R}]}{T_{T}} + P_{C}}, \quad (27)$$

where the expectation is w.r.t. the distribution of H_{SR} . Then the robust average EE maximization problem for the two-hop MIMO relaying network is given by

$$\max_{\boldsymbol{V}_{S}, \boldsymbol{W}_{R}} \min_{\boldsymbol{\Delta}_{RD} \otimes \boldsymbol{E}_{RD}} \operatorname{EE}(\boldsymbol{V}_{S}, \boldsymbol{W}_{R}, \boldsymbol{\Delta}_{RD} \otimes \boldsymbol{E}_{RD}),$$

s.t. $\operatorname{Tr}(\boldsymbol{V}_{S}) \leq P_{S_{\max}},$
 $\operatorname{E}[\operatorname{Tr}(\boldsymbol{W}_{R}(\boldsymbol{H}_{SR}\boldsymbol{V}_{S}\boldsymbol{H}_{SR}^{\mathrm{H}} + \sigma_{r}^{2}\boldsymbol{I}_{N_{R}})\boldsymbol{W}_{R}^{\mathrm{H}})] \leq P_{R_{\max}},$
 $\|\boldsymbol{\Delta}_{RD}\|_{2} \leq \epsilon_{a} \text{ or } \|\boldsymbol{E}_{RD}\|_{2} \leq \epsilon_{m}.$ (28)

By defining the eigenvalue decompositions (EVDs) of R_R and R_S as $R_R = U_R \Lambda_R U_R^H$ and $R_S = U_S \Lambda_S U_S^H$, respectively, where the unitary matrices $U_R \in \mathbb{C}^{N_R \times N_R}$ and $U_S \in \mathbb{C}^{N_S \times N_S}$ consist of the eigenvectors of R_R and R_S , respectively, we have the following Theorem.

Theorem 3: For the robust average EE maximization (28), the worst-case error $\Delta_{RD}^* \otimes E_{RD}^*$ is the same as that given in Theorem 1, with the structures of the optimal V_S^* and W_R^* given by

$$V_{S}^{\star} = U_{S} \Sigma_{S} U_{S}^{\mathrm{H}},$$

$$W_{R}^{\star} = \widehat{Q}_{RD} \Sigma_{X} (I_{N_{R}} + \sigma_{r}^{-2} \Sigma_{SR} \Sigma_{S} \Sigma_{SR}^{\mathrm{H}})^{-\frac{1}{2}} U_{R}^{\mathrm{H}}.$$
 (29)
Proof: See Appendix C.

Based on Theorem 3, the matrix-variable robust average EE problem (28) can be equivalently transformed into a scalar-variable one. As shown in [11], this scalar-variable problem consists of two concave subproblems due to the concavity and monotonicity of the function $E[\log(\cdot)]$. However, evaluating $E[\log(\cdot)]$ imposes high-complexity. To solve (28) efficiently, a deterministic approximation of the average EE is required. We apply Jensen's inequality to the concave function $E[\log(\cdot)]$ to derive the analytical upper bound of the average

$$f_{\rm up}(\boldsymbol{\lambda}_s, \boldsymbol{\sigma}_x) = \frac{\log\left(1 + \frac{\sigma_d^{-2}\sigma_{x,i}^2 \tilde{\sigma}_{rd,i}^2 \sigma_{sr,i}^2 \lambda_{s,i}}{1 + \sigma_r^{-2} \sigma_{sr,i}^2 \lambda_{s,i} + \sigma_r^2 \sigma_d^{-2} \sigma_{x,i}^2 \tilde{\sigma}_{rd,i}^2}\right)}{\sum_{i=1}^{N_S} \lambda_{s,i} + \sum_{i=1}^{N_C} \sigma_r^2 \sigma_{x,i}^2 + P_C} \le \frac{\log\left(1 + \frac{\sigma_d^{-2} \sigma_{x,i}^2 \tilde{\sigma}_{rd,i}^2 \sigma_{sr,i}^2 (\lambda_{s,i} + \sigma_r^2 \sigma_{x,i}^2)}{\sigma_r^{-2} \sigma_{sr,i}^2 + \sigma_r^2 \sigma_d^{-2} \tilde{\sigma}_{rd,i}^2 + 2\sqrt{\sigma_{sr,i}^2 \sigma_{rd,i}^2}}\right)}{\sum_{i=1}^{N_S} \lambda_{s,i} + \sum_{i=1}^{N_C} \sigma_r^2 \sigma_{x,i}^2 + P_C}}.$$
 (25)

EE [11]. Then the proposed alternating optimization can readily be applied to this upper-bound average EE optimization.

Remark 1: The relay needs to feed back the optimal covariance matrix V_S^* to the source, which introduces the feedback errors to V_S^* . When the source covariance matrix error $\Delta V_S \in \mathbb{C}^{N_S \times N_S}$ is taken into account, the proof of Appendix A is not applicable, since both the numerator and denominator of the EE metric contain the semi-infinite ΔV_S . A possible solution is to consider a lower bound optimization of this truly robust EE design, where the possible maximum total power consumption under the spectral norm constrained ΔV_S is adopted. Then ΔV_S is only contained in the rate function and the corresponding minimum achievable rate in (4) is readily observed at $\Delta V_S^* = -\epsilon I_{N_S}$. Since the resultant lowerbound robust EE optimization is similar to that of (8) or (28), the channel-diagonalizing structured new relay beamforming W_R^* can still be proved following the proof in Appendix A.

IV. EXTENSION TO MULTIHOP MIMO AF Relaying Networks

An N_S -antenna source transmits signals to an N_D -antenna destination via K N_R -antenna relays R_k , for $1 \le k \le K$. Denote the source-relay channel by $\boldsymbol{H}_{SR_1} \in \mathbb{C}^{N_R \times N_S}$, the relay-relay channels by $\boldsymbol{H}_{R_k R_{k+1}} \in \mathbb{C}^{N_R \times N_R}$ for $1 \leq 1$ $k \leq K - 1$, and the relay-destination channel by $H_{R_K D} \in$ $\mathbb{C}^{N_D imes N_R}$. The received signals at each relay and the destination are given by (30) and (31), respectively, at the top of the next page. where $\boldsymbol{n}_{R_k} \in \mathbb{C}^{N_R}$, $1 \leq k \leq K$, is the AWGN vector at relay k with the covariance matrix $\sigma_r^2 I_{N_R}$, and $\boldsymbol{n}_{R_D} \in \mathbb{C}^{N_D}$ is the AWGN vector at the destination with the covariance matrix $\sigma_d^2 I_{N_D}$, while W_{R_k} , $1 \le k \le K$, is the beamforming matrix of relay k. Since typically the relays chosen are static or at most slowly mobile w.r.t. the source, the source-relay channel and all relay-relay channels can be acquired with high precision through training [31]. We further assume that the estimated H_{SR_1} and $H_{R_kR_{k+1}}$, $1 \le k \le K-2$, can be transmitted to relay K perfectly. Thus we assume that the perfect $\{H_{SR_1}, H_{R_kR_{k+1}}, 1 \le k \le K-1\}$ are available at relay K, and consider the deterministically imperfect $H_{R_{K}D}$ with the additive or multiplicative CSI errors

$$\boldsymbol{H}_{R_{K}D} = \begin{cases} \widehat{\boldsymbol{H}}_{R_{K}D} + \boldsymbol{\Delta}_{R_{K}D}, & \|\boldsymbol{\Delta}_{R_{K}D}\|_{2} \leq \epsilon_{a}, \\ (\boldsymbol{I}_{N_{D}} + \boldsymbol{E}_{R_{K}D}) \widehat{\boldsymbol{H}}_{R_{K}D}, & \|\boldsymbol{E}_{R_{K}D}\|_{2} \leq \epsilon_{m}, \end{cases} (32)$$

where $\widehat{\boldsymbol{H}}_{R_KD} \in \mathbb{C}^{N_D \times N_R}$ is the known nominal relaydestination channel, $\boldsymbol{\Delta}_{R_KD} \in \mathbb{C}^{N_D \times N_R}$ and $\boldsymbol{E}_{R_KD} \in \mathbb{C}^{N_D \times N_D}$ are the corresponding additive and multiplicative CSI errors. The robust EE optimization problem under the uncertainty model (32) is formulated as

$$\max_{\boldsymbol{V}_{S}, \widetilde{\boldsymbol{W}}_{R}} \min_{\boldsymbol{\Delta}_{R_{K}D} \otimes \boldsymbol{E}_{R_{K}D}} \operatorname{EE}_{M} \left(\boldsymbol{V}_{S}, \widetilde{\boldsymbol{W}}_{R}, \boldsymbol{\Delta}_{R_{K}D} \otimes \boldsymbol{E}_{R_{K}D} \right)$$
$$= \frac{R_{\mathrm{mul}}}{P_{\mathrm{mul}}},$$
s.t. $\operatorname{Tr}(\boldsymbol{V}_{\mathrm{S}}) \leq \operatorname{P}_{\mathrm{S}_{\mathrm{max}}}, \ \operatorname{P}_{\mathrm{R_{k}}} \leq \operatorname{P}_{\mathrm{R}_{\mathrm{max}}}, \ 1 \leq \mathrm{k} \leq \mathrm{K},$
$$\|\boldsymbol{\Delta}_{R_{K}D}\| \leq \epsilon_{a} \text{ or } \|\boldsymbol{E}_{R_{K}D}\| \leq \epsilon_{m},$$
(33)

where $\widetilde{W}_R = \{W_{R_1}, W_{R_2}, \cdots, W_{R_K}\}$, the maximum achievable rate R_{mul} is given by

$$R_{\text{mul}} = \frac{B}{K+1} \log \det \left(\boldsymbol{I}_{N_D} + \boldsymbol{H}_{\text{mul},1} \boldsymbol{H}_{SR_1} \boldsymbol{V}_S \boldsymbol{H}_{SR_1}^{\text{H}} \boldsymbol{H}_{\text{mul},1}^{\text{H}} \boldsymbol{N}_{\text{mul}}^{-1} \right), \quad (34)$$

with

$$\boldsymbol{N}_{\text{mul}} = \sigma_r^2 \sum_{m=1}^{K-1} \boldsymbol{H}_{\text{mul,m}} \boldsymbol{H}_{\text{mul,m}}^{\text{H}} + \sigma_d^2 \boldsymbol{I}_{N_D}, \qquad (35)$$

$$\boldsymbol{H}_{\mathrm{mul},\mathrm{m}} = \boldsymbol{H}_{R_{K}D}\boldsymbol{W}_{R_{K}} \bigg(\prod_{i=m}^{K-1} \boldsymbol{H}_{R_{i}R_{i+1}}\boldsymbol{W}_{R_{i}}\bigg), \quad (36)$$

and the transmit signal power P_{R_k} of relay $k, 1 \le k \le K$, is given by (37) at the top of the next page, while the total power consumption P_{mul} is expressed as

$$P_{\rm mul} = \frac{P_S}{\tau_s} + \sum_{k=1}^{K} \frac{P_{R_k}}{\tau_r} + P'_C,$$
 (38)

in which P'_{C} is the total circuit power consumption. Similarly to (5), $P'_{C} = N_{S}P_{dy,s} + N_{R}\sum_{k=1}^{K}P_{dy,r,k} + P'_{st}$ with $P_{dy,r,k}$ denoting the dynamic power consumption of the *k*th relay's RF chain, $1 \le k \le K$, and the total static power consumption of the source and all relays is $P'_{st} = P_{st,s} + \sum_{k=1}^{K} P_{st,r,k}$.

A. Proposed Robust EE Design

Clearly, the max-min EE problem (33) is more challenging than the problem (8) but the former has the similar structure to the latter and, therefore, it can be solved with the similar approach as detailed in Section III. Specifically, let us define the following SVDs

$$\boldsymbol{H}_{SR_1} = \boldsymbol{U}_{SR_1} \boldsymbol{\Sigma}_{SR_1} \boldsymbol{Q}_{SR_1}^{\mathrm{H}}, \tag{39}$$

$$\boldsymbol{H}_{R_k R_{k+1}} = \boldsymbol{U}_{R_k R_{k+1}} \boldsymbol{\Sigma}_{R_k R_{k+1}} \boldsymbol{Q}_{R_k R_{k+1}}^{\mathsf{H}}, \tag{40}$$

 $\widehat{\boldsymbol{H}}_{R_{K}D} = \widehat{\boldsymbol{U}}_{R_{K}D}\widehat{\boldsymbol{\Sigma}}_{R_{K}D}\widehat{\boldsymbol{Q}}_{R_{K}D}^{\mathrm{H}}, \quad 1 \leq k \leq K-1. \quad (41)$ where $\boldsymbol{U}_{SR_{1}} \in \mathbb{C}^{N_{R} \times N_{R}}$ and $\boldsymbol{Q}_{SR_{1}} \in \mathbb{C}^{N_{S} \times N_{S}}, \quad \boldsymbol{U}_{R_{k}R_{k+1}} \in \mathbb{C}^{N_{R} \times N_{R}}$ and $\boldsymbol{Q}_{R_{k}R_{k+1}} \in \mathbb{C}^{N_{R} \times N_{R}}$, as well as $\widehat{\boldsymbol{U}}_{R_{K}D} \in \mathbb{C}^{N_{D} \times N_{D}}$ and $\widehat{\boldsymbol{Q}}_{R_{K}D} \in \mathbb{C}^{N_{R} \times N_{R}}$ are the unitary matrices for $\boldsymbol{H}_{SR_{1}}, \quad \boldsymbol{H}_{R_{k}R_{k+1}}$ and $\widehat{\boldsymbol{H}}_{R_{K}D}$, respectively, while the diagonal matrices $\boldsymbol{\Sigma}_{SR_{1}} \in \mathbb{C}^{N_{R} \times N_{S}}, \quad \boldsymbol{\Sigma}_{R_{k}R_{k+1}} \in \mathbb{C}^{N_{R} \times N_{R}},$ and $\widehat{\boldsymbol{\Sigma}}_{R_{K}D} \in \mathbb{C}^{N_{D} \times N_{R}}$ contain $N_{P} = \min\{N_{R}, N_{S}\}$ SVs of $\boldsymbol{H}_{SR_{1}}, \quad N_{R}$ SVs of $\boldsymbol{H}_{R_{k}R_{k+1}}$, and $N_{C} = \min\{N_{D}, N_{R}\}$ SVs $\{\widehat{\sigma}_{r_{K}d,1}, \cdots, \widehat{\sigma}_{r_{K}d,N_{C}}\}$ of $\widehat{\boldsymbol{H}}_{R_{K}D}$ at their diagonal positions, respectively.

Theorem 4: For the max-min EE problem (33), the optimal source covariance matrix V_S^* , the optimal relay beamforming matrices $W_{R_k}^*, 1 \leq k \leq K$, and the worst-case errors $\Delta_{R_KD}^* \otimes E_{R_KD}^*$ have the following structures

$$\boldsymbol{V}_{S}^{\star} = \boldsymbol{Q}_{SR_{1}} \boldsymbol{\Sigma}_{S} \boldsymbol{Q}_{SR_{1}}^{\mathrm{H}}, \qquad (42)$$

$$\boldsymbol{W}_{R_{k}}^{\star} = \begin{cases} \boldsymbol{Q}_{R_{1}R_{2}} \boldsymbol{\Sigma}_{\boldsymbol{W}_{R_{1}}} \boldsymbol{U}_{SR_{1}}^{\mathrm{H}}, & k = 1, \\ \boldsymbol{Q}_{R_{k}R_{k+1}} \boldsymbol{\Sigma}_{\boldsymbol{W}_{R_{k}}} \boldsymbol{U}_{R_{k-1}R_{k}}^{\mathrm{H}}, & 2 \leq k \leq K-1, (43) \\ \boldsymbol{\widehat{Q}}_{R_{K}D} \boldsymbol{\Sigma}_{\boldsymbol{W}_{R_{K}}} \boldsymbol{U}_{R_{K-1}R_{K}}^{\mathrm{H}}, & k = K, \end{cases}$$

$$\boldsymbol{\Delta}_{R_{K}D}^{\star} = -\boldsymbol{\widehat{U}}_{R_{K}D} \begin{bmatrix} \boldsymbol{\Lambda}_{R_{K}D} & \boldsymbol{0}_{N_{C} \times (N_{R}-N_{C})} \\ \boldsymbol{0}_{(N_{L}-N_{L}) \times (N_{L}-N_{L})} \boldsymbol{0}_{(N_{L}-N_{L}) \times (N_{L}-N_{L})} \end{bmatrix}$$

$$= -\boldsymbol{U}_{R_{K}D} \begin{bmatrix} \mathbf{0}_{(N_{D}-N_{C})\times N_{C}} & \mathbf{0}_{(N_{D}-N_{C})\times (N_{R}-N_{C})} \end{bmatrix}$$
$$\times \widehat{\boldsymbol{Q}}_{R_{K}D}^{\mathrm{H}} \text{ or } \boldsymbol{E}_{R_{K}D}^{\star} = -\epsilon_{m}\boldsymbol{I}_{N_{D}}, \qquad (44)$$

$$\boldsymbol{y}_{R_{k}} = \begin{cases} \boldsymbol{H}_{SR_{1}}\boldsymbol{s} + \boldsymbol{n}_{R_{1}}, & k = 1, \\ \left(\prod_{i=1}^{k-1} \boldsymbol{H}_{R_{i}R_{i+1}} \boldsymbol{W}_{R_{i}}\right) \boldsymbol{H}_{SR_{1}} \boldsymbol{s} + \sum_{m=1}^{k-1} \prod_{i=m}^{k-1} \boldsymbol{H}_{R_{i}R_{i+1}} \boldsymbol{W}_{R_{i}} \boldsymbol{n}_{R_{i}} + \boldsymbol{n}_{R_{k}}, & k = 2, \cdots, K, \end{cases}$$

$$\boldsymbol{y}_{D} = \boldsymbol{H}_{R_{K}D} \boldsymbol{y}_{R_{K}} + \boldsymbol{n}_{R_{D}} = \boldsymbol{H}_{R_{K}D} \boldsymbol{W}_{R_{K}} \left(\prod_{i=1}^{K-1} \boldsymbol{H}_{R_{i}R_{i+1}} \boldsymbol{W}_{R_{i}}\right) \boldsymbol{H}_{SR_{1}} \boldsymbol{s} \\ + \boldsymbol{H}_{R_{K}D} \boldsymbol{W}_{R_{K}} \left(\sum_{m=1}^{K-1} \prod_{i=m}^{K-1} \boldsymbol{H}_{R_{i}R_{i+1}} \boldsymbol{W}_{R_{i}} \boldsymbol{n}_{R_{i}}\right) + \boldsymbol{H}_{R_{K}D} \boldsymbol{W}_{R_{K}} \boldsymbol{n}_{R_{K}} + \boldsymbol{n}_{R_{D}}, \qquad (31)$$

$$P_{R_{k}} = \begin{cases} \operatorname{Tr}\left(\boldsymbol{W}_{R_{1}}\left(\boldsymbol{H}_{SR_{1}}\boldsymbol{V}_{S}\boldsymbol{H}_{SR_{1}}^{\mathrm{H}} + \sigma_{r}^{2}\boldsymbol{I}_{N_{R}}\right)\boldsymbol{W}_{R_{1}}^{\mathrm{H}}\right), & k = 1, \\ \operatorname{Tr}\left(\boldsymbol{W}_{R_{k}}\left(\prod_{i=1}^{k-1}\boldsymbol{H}_{R_{i}R_{i+1}}\boldsymbol{W}_{R_{i}}\right)\boldsymbol{H}_{SR_{1}}\boldsymbol{V}_{S}\boldsymbol{H}_{SR_{1}}^{\mathrm{H}}\left(\prod_{i=1}^{k-1}\boldsymbol{H}_{R_{i}R_{i+1}}\boldsymbol{W}_{R_{i}}\right)^{\mathrm{H}} \\ + \sigma_{r}^{2}\sum_{m=1}^{k-1}\left(\prod_{i=m}^{k-1}\boldsymbol{H}_{R_{i}R_{i+1}}\boldsymbol{W}_{R_{i}}\right)\left(\prod_{i=m}^{k-1}\boldsymbol{H}_{R_{i}R_{i+1}}\boldsymbol{W}_{R_{i}}\right)^{\mathrm{H}} + \sigma_{r}^{2}\boldsymbol{I}_{N_{R}}\right)\boldsymbol{W}_{R_{k}}^{\mathrm{H}}\right), \quad 2 \leq k \leq K, \end{cases}$$
(37)

where $\Lambda_{R_K D} = \text{diag} \{\min\{\widehat{\sigma}_{r_K d, 1}, \epsilon_a\}, \cdots, \min\{\widehat{\sigma}_{r_K d, N_C}, \epsilon_a\}\},\$

$$\boldsymbol{\Sigma}_{\boldsymbol{W}_{R_1}} = \boldsymbol{\Sigma}_{X_1} \left(\boldsymbol{I}_{N_R} + \sigma_r^{-2} \boldsymbol{\Sigma}_{SR_1} \boldsymbol{\Sigma}_S \boldsymbol{\Sigma}_{SR_1}^{\mathrm{H}} \right)^{-\frac{1}{2}}, \tag{45}$$

$$\Sigma_{\boldsymbol{W}_{R_{k}}} = \Sigma_{X_{k}} \Big(\boldsymbol{I}_{N_{R}} + \sum_{m=1}^{k-1} \overline{\widetilde{\boldsymbol{\Sigma}}}_{k}^{m} + \sigma_{r}^{-2} \Big(\prod_{i=1}^{k-1} \widetilde{\boldsymbol{\Sigma}}_{i,i+1} \Big) \Sigma_{SR_{1}} \Sigma_{S} \\ \times \cdot \Sigma_{SR_{1}}^{\mathrm{H}} \Big(\prod_{i=1}^{k-1} \widetilde{\boldsymbol{\Sigma}}_{i,i+1} \Big)^{\mathrm{H}} \Big)^{-\frac{1}{2}}, 2 \leq k \leq K,$$
(46)

with
$$\Sigma_{X_k} = \operatorname{diag} \{ \sigma_{x_k,1}, \cdots, \sigma_{x_k,N_C}, \underbrace{0,0,\cdots,0}_{N_R-N_C} \}$$
 for $1 \leq i \leq K$, $\overline{\widetilde{\Sigma}}_k^m = (\prod_{i=m}^{k-1} \widetilde{\Sigma}_{i,i+1}) (\prod_{i=m}^{k-1} \widetilde{\Sigma}_{i,i+1})^{\mathrm{H}}$ and $\widetilde{\Sigma}_{i,i+1} = \sum_{i=m}^{k-1} \widetilde{\Sigma}_{i,i+1} = \sum_{i=m}^{k-1} \widetilde{\Sigma}_{i,i$

 $\mathbf{\Sigma}_{R_iR_{i+1}}\mathbf{\Sigma}_{\mathbf{W}_{R_i}}.$

Proof: See Appendix D. \Box

Similarly to the two-hop case, based on Theorem 4 and Lemma 1, we can equivalently transform the robust EE problem (33) into the follow optimization problem with scalar variables

where $\tilde{\Sigma}_{R_KD} = (\hat{\Sigma}_{R_KD} - \Lambda_{R_KD}) \Sigma_{W_{R_K}}$ and $\tilde{\Sigma}_{R_KD} = (1-\epsilon_m) \hat{\Sigma}_{R_KD} \Sigma_{W_{R_K}}$ are defined for the additive and multiplicative CSI errors, respectively, and η_{mul} is the auxiliary variable. For convenience, denote $\Sigma_S = \Sigma_{X_0}$. Observe that the problem (47) is convex w.r.t Σ_{X_k} when the remaining variables $\{\Sigma_{X_0}, \Sigma_{X_1}, \dots, \Sigma_{X_K}\} \setminus \Sigma_{X_k}$ are fixed. Therefore, we can apply the alternating optimization of Section III-B to efficiently solve the problem (47) by decomposing it into (K+1) alternating subproblems, in order to obtain a locally optimal solution.

B. Extension to Imperfect H_{SR_1} and $H_{R_kR_{k+1}}$, $\forall k$

We now consider the most generic case, where the source-relay channel H_{SR_1} and all the relay-relay channels $H_{R_kR_{k+1}}$, $\forall k$, are also imperfect at relay R_K . Similarly to the two-hop case, by adopting the statistically imperfect

 H_{SR_1} and $H_{R_kR_{k+1}}$, $\forall k$, and the deterministically imperfect $H_{R_{K}D}$, the robust average EE optimization for this generic multihop AF relaying network can be formed. Following the same philosophy in proving Theorem 3, a similar conclusion can also be obtained for the multihop scenario by proving the optimal channel-diagonalizing structure one by one for the optimal source covariance matrix V_S^{\star} and the optimal relay beamforming matrices $\boldsymbol{W}^*_{R_k}$, $1 \le k \le K$. Specifically, the eigenspaces of V_S^* and $W_{R_k}^*$, $1 \le k \le K-1$, are aligned with that of the source spatial correlation matrix and that of the relay R_k spatial correlation matrix, respectively, while $W^*_{R_K}$ is jointly determined by the eigenspace of the relay R_K spatial correlation matrix and the right singular matrix of the relay-destination channel H_{R_KD} . The corresponding robust average EE problem can also be transformed into a robust optimization with scalar variables. However, due to the simultaneous expectations for multiple statistically imperfect channels, this scalar-variable problem is generally intractable and cannot be decomposed into a series of convex subproblems [13]. A possible solution is to apply successive Jensen's inequalities to the expectations on $\{H_{SR_1}, H_{R_k, R_{k+1}}, \forall k\}$ to find an upper-bound of the average EE [13], which can then be solved by the proposed alternating optimization.

Remark 2: The last relay R_K needs to feed back the optimal source covariance matrix V_S^* and optimal relay beamforming matrices $W_{R_k}^*$, $1 \le k \le K - 1$, perfectly to the source and corresponding relays. If the feedback errors for $\{V_S^*, W_{R_k}^*, 1 \le k \le K - 1\}$ are serious, they should be additionally imposed on the robust EE design, and the optimal beamforming matrix $W_{R_K}^*$ of relay K should also be redesigned accordingly. Unfortunately, even if the spectral norm constrained errors for $\{V_S^*, W_{R_k}^*, 1 \le k \le K - 1\}$ are jointly considered, the channel-diagonalizing structured optimal $W_{R_K}^*$ is not guaranteed, since the multiple relay beamforming errors are coupled in both the objective and the transmit power constraints. Future research is warranted to develop the low-complexity suboptimal algorithms to effectively address this issue.



Fig. 2. Convergence of Algorithm 1: (a) inner Dinkelbach loop of step 2 at the first outer iteration, (b) outer alternating optimization loop, given two sets of different initial points and for both additive and multiplicative CSI errors, and (c) inner Dinkelbach loop of step 2 at the first outer iteration with error bars, for one set of initial points and additive CSI errors. $N_S/N_R/N_D = 4/6/4$.

V. SIMULATION STUDY

In the simulation, the source is a base station (BS), while the destination and relays are mobile stations (MSs). The default parameters of the simulated MIMO AF relaying network are listed in Table II. Unless otherwise stated, these default values are used. To demonstrate the excellent performance of our robust EE design, we adopt the non-robust EE maximization (NREE) and the naive AF based EE maximization (NAF) [17] for comparison. For the NREE scheme, the optimization problem (8)/(32) is firstly solved by assuming no CSI errors, i.e., $\epsilon_a = \epsilon_m = 0$. Then the resultant optimal solution is applied to the imperfect CSI scenario for calculating the worst-case EE. For the NAF scheme [17], the relay scales the received signal transmitted by source with the maximum power by a constant to realize the maximum relay power transfer. All simulation results are obtained by averaging over 100 channel realizations.

A. Two-Hop MIMO AF Relaying Networks

The convergence of the proposed alternative optimization algorithm is investigated under the two sets of the initial values $\{\lambda_s^{(0)}, \eta^{(0)}\}$, given by $\{\lambda_{s,\text{ini1}}^{(0)}, \eta_1^{(0)}\} = \{\frac{P_{S_{\text{max}}}}{2N_S} \mathbf{1}_{N_S}, 0\}$ and $\{\lambda_{s,\text{ini2}}^{(0)}, \eta_2^{(0)}\} = \{\frac{P_{S_{\text{max}}}}{2N_S} [\mathbf{1}_{N_S/2}^T \mathbf{0}_{N_S/2}^T]^T, 0.1\}$. Algorithm 1 consists of an outer alternating optimization loop, and within each alternating iteration, there are two inner Dinkelbach iterative loops at Step 2 and step 3. To demonstrate the convergence of the two inner Dinkelbach iterative loops, Fig. 2 (a) plots the objective value of the problem (20) after each Dinkelbach

² In our simulations, the large-scale fading coefficient σ_h^2 is determined according to the LTE path-loss model $-\sigma_h^2 + 30 = L_P(dB) = 128.1 + 37.6 \log(d \text{ [Km]})$ with d = 330 m [33].

iteration for optimizing σ_x given λ_s at the first outer iteration. Observe from Fig. 2 (a) that this very first Dinkelbach iterative procedure of Algorithm 1 takes no more than 6 iterations to converge. Since any subsequent Dinkelbach iterative procedure is unlikely to take more iterations to converge, we conclude that for this example, any Dinkelbach iterative procedure of Algorithm 1 takes no more than 6 iterations to converge. The convergence of Algorithm 1 is illustrated in Fig. 2 (b), where it is seen that for this example, Algorithm 1 takes no more than 7 outer iterations to converge. Fig. 2 (c) depicts the curve of Fig. 2 (a) corresponding to the additive CSI errors with the initial condition $\{\lambda_{s,ini1}^{(0)}, \eta_1^{(0)}\}$ but with the logarithmic scale in *y*-axis and with the error bars. It can be seen that after 6 iterations, the objective value becomes smaller than 10^{-3} .

The influence of antenna configuration to the worst-case EE performance is investigated in Fig. 3, which plots the worst-case EE performance of the proposed robust EE design as the functions of the source maximum transmit power $P_{S_{\text{max}}}$ under different antenna configurations $N_S/N_R/N_D$. We observe that the antenna configuration 4/4/6 attains the highest worst-case EE, and the antenna configuration 4/6/4attains the second highest worst-case EE, which is considerably larger than the 4/4/4 configuration. While the 6/4/4configuration achieves the lowest worst-case EE. We now explain these phenomena based on the following facts. First, the circuit power consumption at receiver is neglected in our work, since reception generally consumes much less circuit power than transmission [23]. Second, the existing literature [7], [34] have shown that increasing the number of source antennas N_S without optimizing the source covariance matrix generally causes the capacity shrinking phenomenon for MIMO AF relaying systems. Even considering the design

$$\max_{\boldsymbol{\Sigma}_{S},\forall\boldsymbol{\Sigma}_{X_{k}}} \log \det \left(\boldsymbol{I}_{N_{D}} + \left(\widetilde{\boldsymbol{\Sigma}}_{R_{K}D} \left(\prod_{i=1}^{K-1} \widetilde{\boldsymbol{\Sigma}}_{i,i+1} \right) \boldsymbol{\Sigma}_{SR_{1}} \boldsymbol{\Sigma}_{S} \boldsymbol{\Sigma}_{SR_{1}}^{\mathrm{H}} \left(\prod_{i=1}^{K-1} \widetilde{\boldsymbol{\Sigma}}_{i,i+1} \right)^{\mathrm{H}} \widetilde{\boldsymbol{\Sigma}}_{R_{K}D}^{\mathrm{H}} \right) \\ \times \left(\sigma_{r}^{2} \sum_{m=1}^{K-1} \widetilde{\boldsymbol{\Sigma}}_{R_{K}D} \overline{\widetilde{\boldsymbol{\Sigma}}}_{K}^{m} \widetilde{\boldsymbol{\Sigma}}_{R_{K}D}^{\mathrm{H}} + \sigma_{d}^{2} \boldsymbol{I}_{N_{D}} \right)^{-1} \right) - \eta_{\mathrm{mul}} \left(\mathrm{Tr}(\boldsymbol{\Sigma}_{S}) + \sum_{k=1}^{K} \mathrm{Tr}(\boldsymbol{\Sigma}_{X_{k}}) + P_{C} \right),$$

s.t. $\mathrm{Tr}(\boldsymbol{\Sigma}_{S}) \leq P_{S_{\mathrm{max}}}, \ \mathrm{Tr}(\boldsymbol{\Sigma}_{X_{k}}) \leq P_{R_{k_{\mathrm{max}}}}, \ 1 \leq k \leq K,$ (47)

TABLE II					
DEFAULT SYSTEM PARAMETERS					

Parameter	Value	Parameter	Value	
The number of BS/Relay/MS antennas	4/6/4 (4/4/4;	Relay dynamic power consumption $P_{dy,r}$	35dBm [32]	
$N_S/N_R/N_D$	6/4/4; 4/4/6)			
Transmission bandwidth B	50MHz	Relay static power consumption $P_{st,r}$	30dBm [32]	
Power amplifier efficiency τ_s, τ_r	0.5	Additive errors threshold (two-hop) ϵ_a	$\epsilon_a = p \ \widehat{H}_{RD} \ _2, p = 0.4$	
BS maximum transmit power $P_{S_{\text{max}}}$	46dBm [33]	Additive errors threshold (multihop) ϵ_a	$\epsilon_a = p \ \widehat{H}_{R_K D} \ _2, p = 0.4$	
Relay maximum transmit power $P_{R_{\text{max}}}$	40dBm [33]	Multiplicative errors threshold ϵ_m	$\epsilon_m = p = 0.4$	
BS dynamic power consumption $P_{dy,s}$	40dBm [32]	Rayleigh MIMO relaying channels with	$\mathcal{CN}(0,\sigma_h^2 \boldsymbol{I}), \sigma_h^2 =$	
		large-scale fading coefficient σ_h^2 . ²	-80dBm [33]	
BS static power consumption $P_{st,s}$	35dBm [32]	Noise power $\sigma_r^2 = \sigma_d^2 = N_0 B$	-90 dBm (N ₀ =	
			-167dBm/Hz) [32]	



Fig. 3. Comparison of the worst-case EE performance as the functions of the source maximum transmit power $P_{S_{\max}}$ given different antenna configurations $N_S/N_R/N_D$.

of source covariance matrix, the capacity gain is close to zero as N_S increases [7], [34]. By contrast, adding more relay and/or destination antennas is helpful to improve system capacity, of which the enhancement is more evident when increasing the number of relay antennas [34].

These known conclusions explain why the best worst-case EE performance for both multiplicative and additive CSI errors is observed at $N_S/N_R/N_D = 4/4/6$, which is because the achievable data rate increases with the number of destination antennas N_D , while the transmit power consumption remains unchanged. These conclusions also agree with our observation that the achievable worst-case EE under $N_S/N_R/N_D = 4/6/4$ is significantly higher than that under $N_S/N_R/N_D = 4/4/4$, because when the number of relay antennas N_R increases, the remarkable increase of data rate outweighs the increase in dynamic relay circuit power consumption. However, for the case of $N_S/N_R/N_D = 6/4/4$, the achievable worst-case EE is much reduced, since the increased source dynamic circuit power consumption outweighs the slight data rate gain due to the increasing number of source antennas N_S . Our results therefore support the existing literature and provide insights to design MIMO AF relaying systems, namely, increasing the number of relay and/or destination antennas rather than the number of source antennas is beneficial to improve system's worst-case EE performance.



Fig. 4. Comparison of the worst-case EE performance as the functions of: (a) the source maximum transmit power $P_{S_{\text{max}}}$, and (b) the relay maximum transmit power $P_{R_{\text{max}}}$, for three designs. $N_S/N_R/N_D = 4/6/4$.

Fig. 4(a) compares the worst-case EE performance as the functions of $P_{S_{\max}}$ for the proposed robust EE design, NREE and NAF schemes, while Fig. 4(b) shows the worst-case EE performance as the functions of $P_{R_{\text{max}}}$ for the three designs. As expected, the proposed robust EE design achieves the highest worst-case EE, while the NAF scheme has the worst-case EE. For our robust EE design and NREE scheme, the worst-case EE first increases with $P_{S_{\text{max}}}$ ($P_{R_{\text{max}}}$) but becomes saturated for large $P_{S_{\text{max}}}$ ($P_{R_{\text{max}}}$). This is because when $P_{S_{\text{max}}}$ ($P_{R_{\text{max}}}$) is small compared to the circuit power consumption P_C , the EE metric is mainly determined by the maximum achievable rate in the numerator of the EE metric. Increasing $P_{S_{\text{max}}}(P_{R_{\text{max}}})$ increases the source (relay) transmit power too, which in turn increases the maximum achievable rate. Therefore, for relatively small $P_{S_{\text{max}}}$ ($P_{R_{\text{max}}}$), increasing $P_{S_{\text{max}}}$ ($P_{R_{\text{max}}}$) increases the worst-case EE. In this region, the power constraint is active, i.e., the source (relay) transmit power reaches $P_{S_{\max}}$ ($P_{R_{\max}}$). By contrast, when $P_{S_{\max}}$ ($P_{R_{\max}}$) becomes large, the EE metric is also determined by the source (relay) transmit power in the denominator of the EE metric. Therefore, to maximize the worst-case EE, a trade-off between the maximum achievable rate and the total transmit power must be made. As a result, the maximum worst-case EE metric reaches a saturated value. In this region, the power constraint is inactive and the source (relay) transmit power remains constant. For the NAF, the worst-case EE decreases with $P_{S_{\text{max}}}$ ($P_{R_{\text{max}}}$) for large $P_{S_{\text{max}}}$ ($P_{R_{\text{max}}}$). This



Fig. 5. Influence of CSI uncertainty p on the achievable worst-case EE for three designs. $N_S/N_R/N_D\,{=}\,4/6/4.$

is because the maximum achievable rate of the NAF strategy cannot be arbitrarily enhanced by increasing $P_{S_{\text{max}}}$ ($P_{R_{\text{max}}}$), since the relay simultaneously amplifies the received signal and noise. Therefore, for large $P_{S_{\text{max}}}$ ($P_{R_{\text{max}}}$), the NAF's worst-case EE decreases with $P_{S_{\text{max}}}$ ($P_{R_{\text{max}}}$) due to the limited maximum achievable rate and the increase in the total power consumption.

The two thresholds are fairly set for the multiplicative and additive CSI errors in Table II, as they correspond to the same quantitative measure for the spectral norm constrained multiplicative and additive CSI errors. According to Theorem 1, we can infer that the SVs of the worst-case relay-destination channel for multiplicative CSI errors are larger or equal to those for additive CSI errors. Therefore, we can conclude that under the same size of the spectral norm constrained CSI errors, the quality of the relay-destination channel under multiplicative CSI errors is better than that under additive errors. Naturally, the achievable worst-case EE under multiplicative CSI errors is also higher than that under additive CSI errors. This is also confirmed by Fig. 4.

Fig. 5 investigates the impact of the CSI uncertainty threshold p on the achievable worst-case EE. Not surprisingly, as pincreases, the worst-case EE performance decreases for every scheme. Again our robust design achieves the best worst-case EE, while the NAF has the worst performance. Moreover, the performance gap between our robust EE design and the NREE increases with p, which further indicates the effectiveness of our design. Fig. 6 shows the worst-case EE performance as the functions of the circuit power consumption P_C for three schemes. As expected, our robust design attains the highest worst-case EE. Observe that the worst-case EE decreases as P_c increases, since P_c only appears in the denominator of the EE metric. Next we consider the statistically imperfect source-relay channel H_{SR} . The source and relay spatial correlation matrices R_R and R_S are simulated using the exponential model [35]. Specifically, for $i, j = 1, \dots, N_R(N_S)$ and $j \geq i$, the (i, j)-th elements of \boldsymbol{R}_R and \boldsymbol{R}_S are given respectively by $[\mathbf{R}_R]_{i,j} = p_r^{j-i}$ and $[\mathbf{R}_S]_{i,j} = p_s^{j-i}$, where $|p_r| < 1$ and $|p_s| < 1$ [35]. Without loss of generality, only additive CSI error is considered for the relay-destination channel. Fig. 7 depicts the worst-case EE performance versus $P_{S_{\text{max}}}$ for the



Fig. 6. Influence of circuit power consumption P_C on the achievable worstcase EE for three designs. $N_S/N_R/N_D = 4/6/4$.



Fig. 7. Influence of the source maximum transmit power $P_{S_{\rm max}}$ on the achievable worst-case EE for two designs. The statistically imperfect source-relay channel and additive CSI errors for the relay-destination channel are considered. $N_S/N_R/N_D \!=\! 4/6/4.$

upper-bound robust average EE design of Section III-E and the NREE scheme. Observe from Fig. 7 that the relationship between the worst-case EE and $P_{S_{\text{max}}}$ for the two designs is similar to that shown in Fig. 4 (a). Compared with the case of $p_s = p_r = 0.3$, the stronger correlation of $p_s = p_r = 0.5$ leads to higher achievable worst-case EE, because higher correlation means higher source-relay channel energy, which is beneficial to improve the ergodic rate. For the two designs considered, Fig. 8 shows that the achievable worst-case EE decreases with the BS dynamic power consumption $P_{dy,s}$, which is a portion of the total circuit power consumption P_C . The reason is similar to that given for Fig. 6.

B. Three-Hop (K = 2) MIMO AF Relaying Networks

A three-hop MIMO AF relaying network is simulated, and we only consider additive CSI errors for the relay-destination channel. Fig. 9 (a) compares the worst-case EE performance versus the relays' maximum transmit power $P_{R_{\text{max}}}$ for the three designs. Compared with Fig. 4 (b), similar trends between the worst-case EE and $P_{R_{\text{max}}}$ for the three schemes can also be observed from Fig. 9 (a). Clearly, the achievable worst-case EE in the three-hop case is lower than that in the



Fig. 8. Influence of the BS dynamic power consumption $P_{dy,s}$ on the achievable worst-case EE for two designs. The statistically imperfect source-relay channel and additive CSI errors for the relay-destination channel are considered. $N_S/N_R/N_D = 4/6/4$.



Fig. 9. Comparison of the worst-case EE performance as the functions of: (a) the relays' maximum transmit power $P_{R_{\max}}$, and (b) the CSI uncertainty p, for three designs. The three-hop MIMO AF relaying network under additive CSI uncertainty is considered. $N_S = 4$, $N_{R_1} = N_{R_2} = 6$ and $N_D = 4$.

two-hop case due to the greater channel fading and higher power consumption. Fig. 9(b) depicts the influence of p on the worst-case EE for the three designs. Both Fig. 9(a) and Fig. 9(b) confirm that the proposed robust EE design attains the best worst-case EE performance.

VI. CONCLUSIONS

We have optimized the EE of two-hop MIMO AF relaying networks under the deterministically imperfect CSI. By considering the additive and multiplicative CSI errors for the relay-destination channel, the source covariance and relay beamforming matrices are jointly optimized to maximize the worst-case EE. We have proved the existence of a saddle point for this robust EE problem, and have derived the channel-diagonalizing structure of the optimal source covariance and relay beamforming matrices as well as the worst-case errors under the spectral-norm constrained additive and multiplicative CSI errors. Based on this structure, the original robust EE problem is transformed into an optimization problem with scalar variables, which can be efficiently solved by the proposed alternating optimization. We have also proved that all these results are applicable when the statistically imperfect source-relay channel is additionally imposed. Furthermore,

we have extended our work to multihop MIMO AF relaying networks, and have proved that the channel-diagonalizing structure remains optimal for the source covariance matrix and all the relays' beamforming matrices under deterministically imperfect relay-destination CSI.

Appendix

A. Proof of Theorem 1

Proof: Unless otherwise stated, the eigenvalues (EVs)/SVs of an EVD/SVD for a matrix are always arranged in a decreasing order. First we have the following lemma [36].c

Lemma 2: For the two $N \times N$ Hermitian matrices A and B whose EVs are denoted by $\lambda_i(A)$ and $\lambda_i(B)$, respectively, for $i = 1, \dots, N$, we have

$$\begin{split} \prod_{i=1}^{N} \left(\lambda_{i}(\boldsymbol{A}) + \lambda_{i}(\boldsymbol{B}) \right) &\leq \det(\boldsymbol{A} + \boldsymbol{B}) \\ &\leq \prod_{i=1}^{N} \left(\lambda_{i}(\boldsymbol{A}) + \lambda_{N+1-i}(\boldsymbol{B}) \right), \end{split} \tag{48}$$

$$\sum_{i=1}^{N} \lambda_{i}(\boldsymbol{A}) \lambda_{N+1-i}(\boldsymbol{B}) &\leq \operatorname{Tr}(\boldsymbol{A}\boldsymbol{B}) \leq \sum_{i=1}^{N} \lambda_{i}(\boldsymbol{A}) \lambda_{i}(\boldsymbol{B}). \end{aligned} \tag{49}$$

All the equalities in (48) and (49) hold only when A and B are simultaneously diagonalizable.

1) Optimal W_{R}^{\star} : Re-express the rate formulation (9) as

According to the identity $\det(I_N + AB) = \det(I_K + BA)$, where $A \in \mathbb{C}^{N \times K}$ and $B \in \mathbb{C}^{K \times N}$, the achievable EE metric $\operatorname{EE}(V_S, W_R, \Delta_{RD} \otimes E_{RD})$ can be reformulated as (50), shown at the top of the next page.

Next perform the EVDs of the nonnegative definite matrices $H_{SR}V_SH_{SR}^{H}$ and $H_{RD}^{H}H_{RD}$:

$$\boldsymbol{H}_{SR}\boldsymbol{V}_{S}\boldsymbol{H}_{SR}^{\mathrm{H}} = \widetilde{\boldsymbol{U}}_{SR}\widetilde{\boldsymbol{\Sigma}}_{SR}\widetilde{\boldsymbol{U}}_{SR}^{\mathrm{H}}, \qquad (52)$$

$$\boldsymbol{H}_{RD}^{\mathrm{H}}\boldsymbol{H}_{RD} = \widetilde{\boldsymbol{Q}}_{RD}\widetilde{\boldsymbol{\Sigma}}_{RD}\widetilde{\boldsymbol{Q}}_{RD}^{\mathrm{H}}, \qquad (53)$$

where $\widetilde{\Sigma}_{SR} = \text{diag}\{\widetilde{\sigma}_{sr,1}^2, \cdots, \widetilde{\sigma}_{sr,N_P}^2, 0, \cdots, 0\}$ and $\widetilde{\Sigma}_{RD} = \text{diag}\{\widetilde{\sigma}_{rd,1}^2, \cdots, \widetilde{\sigma}_{rd,N_C}^2, 0, \cdots, 0\}$ contain the N_P nonzero EVs of $H_{SR}V_SH_{SR}^H$ and the N_C nonzero EVs of $H_{RD}^HH_{RD}$, respectively, while $\widetilde{U}_{SR} \in \mathbb{C}^{N_R \times N_R}$ and $\widetilde{Q}_{RD} \in \mathbb{C}^{N_R \times N_R}$ are the associated unitary matrices. Note that \widetilde{Q}_{RD} and $\widetilde{\Sigma}_{RD}$ are unknown, while $\widetilde{\Sigma}_{SR}$ depends on the matrix variable V_S whose optimal structure is yet to be determined. Clearly, $\widetilde{\Sigma}_{SR} \succeq \mathbf{0}$ and $\widetilde{\Sigma}_{RD} \succeq \mathbf{0}$. By defining $\mathbf{X} \in \mathbb{C}^{N_R \times N_R}$ as

$$\boldsymbol{X} = \widetilde{\boldsymbol{Q}}_{RD}^{\mathrm{H}} \boldsymbol{W}_{R} \widetilde{\boldsymbol{U}}_{SR} \big(\boldsymbol{I}_{N_{R}} + \sigma_{r}^{-2} \widetilde{\boldsymbol{\Sigma}}_{SR} \big)^{\frac{1}{2}}, \qquad (54)$$

we can express the relay beamforming matrix W_R as

$$\boldsymbol{W}_{R} = \widetilde{\boldsymbol{Q}}_{RD} \boldsymbol{X} \left(\boldsymbol{I}_{N_{R}} + \sigma_{r}^{-2} \widetilde{\boldsymbol{\Sigma}}_{SR} \right)^{-\frac{1}{2}} \widetilde{\boldsymbol{U}}_{SR}^{\mathrm{H}}.$$
(55)

Then X is the new optimization matrix variable. Substituting (52), (53) and (55) into (51) (at the top of the next page) yields (56) (also at the top of the next page). Denote $X^{H} \widetilde{\Sigma}_{RD} X = U_{T} \Sigma_{T} U_{T}^{H}$, where the unitary matrix $U_{T} \in \mathbb{C}^{N_{R} \times N_{R}}$ and the $N_{R} \times N_{R}$ diagonal matrix Σ_{T} contains N_{C} nonzero EVs of $X^{H} \widetilde{\Sigma}_{RD} X$. For any $X^{H} \widetilde{\Sigma}_{RD} X$,

$$2^{\frac{2R_D}{B}} = \frac{\det \left(\boldsymbol{I}_{N_R} + \sigma_r^2 \sigma_d^{-2} \left(\boldsymbol{I}_{N_R} + \sigma_r^{-2} \boldsymbol{H}_{SR} \boldsymbol{V}_S \boldsymbol{H}_{SR}^{\mathrm{H}} \right) \boldsymbol{W}_R^{\mathrm{H}} \boldsymbol{H}_{RD}^{\mathrm{H}} \boldsymbol{H}_{RD} \boldsymbol{W}_R \right)}{\det \left(\boldsymbol{I}_{N_R} + \sigma_r^2 \sigma_d^{-2} \boldsymbol{W}_R^{\mathrm{H}} \boldsymbol{H}_{RD}^{\mathrm{H}} \boldsymbol{H}_{RD} \boldsymbol{W}_R \right)} \qquad (50)$$

$$EE \left(\boldsymbol{V}_S, \boldsymbol{W}_R, \boldsymbol{\Delta}_{RD} \varnothing \boldsymbol{E}_{RD} \right) = \frac{\frac{B}{2} \log \det \left(\boldsymbol{I}_{N_D} + \sigma_r^2 \sigma_d^{-2} \boldsymbol{H}_{RD} \boldsymbol{W}_R \left(\boldsymbol{I}_{N_R} + \sigma_r^{-2} \boldsymbol{H}_{SR} \boldsymbol{V}_S \boldsymbol{H}_{SR}^{\mathrm{H}} \right) \boldsymbol{W}_R^{\mathrm{H}} \boldsymbol{H}_{RD}^{\mathrm{H}} \right)}{\operatorname{Tr} \left(\boldsymbol{V}_S \right) + \operatorname{Tr} \left(\boldsymbol{W}_R \left(\boldsymbol{H}_{SR} \boldsymbol{V}_S \boldsymbol{H}_{SR}^{\mathrm{H}} + \sigma_r^2 \boldsymbol{I}_{N_R} \right) \boldsymbol{W}_R^{\mathrm{H}} \right) + P_C} - \frac{\frac{B}{2} \log \det \left(\boldsymbol{I}_{N_R} + \sigma_r^2 \sigma_d^{-2} \boldsymbol{W}_R^{\mathrm{H}} \boldsymbol{H}_{RD}^{\mathrm{H}} \boldsymbol{H}_{RD} \boldsymbol{W}_R \right)}{\operatorname{Tr} \left(\boldsymbol{V}_S \right) + \operatorname{Tr} \left(\boldsymbol{W}_R \left(\boldsymbol{H}_{SR} \boldsymbol{V}_S \boldsymbol{H}_{SR}^{\mathrm{H}} + \sigma_r^2 \boldsymbol{I}_{N_R} \right) \boldsymbol{W}_R^{\mathrm{H}} \right) + P_C}$$

$$\frac{2}{B} \cdot \text{EE}(\boldsymbol{V}_{S}, \boldsymbol{X}, \boldsymbol{\Delta}_{RD} \otimes \boldsymbol{E}_{RD}) = \frac{\log \det \left(\boldsymbol{I}_{N_{R}} + \sigma_{r}^{2} \sigma_{d}^{-2} \boldsymbol{X}^{\mathrm{H}} \widetilde{\boldsymbol{\Sigma}}_{RD} \boldsymbol{X}\right)}{\operatorname{Tr}(\boldsymbol{V}_{S}) + \operatorname{Tr}(\sigma_{r}^{2} \boldsymbol{X} \boldsymbol{X}^{\mathrm{H}}) + P_{C}} + \frac{\log \det \left(\boldsymbol{I}_{N_{R}} + \sigma_{r}^{-2} \widetilde{\boldsymbol{\Sigma}}_{SR}\right)}{\operatorname{Tr}(\boldsymbol{V}_{S}) + \operatorname{Tr}(\sigma_{r}^{2} \boldsymbol{X} \boldsymbol{X}^{\mathrm{H}}) + P_{C}} - \frac{\log \det \left(\boldsymbol{I}_{N_{R}} + \sigma_{r}^{-2} \widetilde{\boldsymbol{\Sigma}}_{SR} + \sigma_{r}^{2} \sigma_{d}^{-2} \boldsymbol{X}^{\mathrm{H}} \widetilde{\boldsymbol{\Sigma}}_{RD} \boldsymbol{X}\right)}{\operatorname{Tr}(\boldsymbol{V}_{S}) + \operatorname{Tr}(\sigma_{r}^{2} \boldsymbol{X} \boldsymbol{X}^{\mathrm{H}}) + P_{C}} - \frac{\log \det \left(\boldsymbol{I}_{N_{R}} + \sigma_{r}^{-2} \widetilde{\boldsymbol{\Sigma}}_{SR} + \sigma_{r}^{2} \sigma_{d}^{-2} \boldsymbol{X}^{\mathrm{H}} \widetilde{\boldsymbol{\Sigma}}_{RD} \boldsymbol{X}\right)}{\operatorname{Tr}(\boldsymbol{V}_{S}) + \operatorname{Tr}(\sigma_{r}^{2} \boldsymbol{X} \boldsymbol{X}^{\mathrm{H}}) + P_{C}}$$
(56)

by introducing $\widetilde{X} = XU_T$, we have $\widetilde{X}^{H} \widetilde{\Sigma}_{RD} \widetilde{X} = \Sigma_T$, i.e., $\widetilde{X}^{H} \widetilde{\Sigma}_{RD} \widetilde{X}$ is diagonal, $\operatorname{Tr}(XX^{H}) = \operatorname{Tr}(\widetilde{X}\widetilde{X}^{H})$ and $\det (I_{N_R} + \sigma_r^2 \sigma_d^{-2} X^{H} \widetilde{\Sigma}_{RD} X)$ $= \det (I_{N_R} + \sigma_r^2 \sigma_d^{-2} \widetilde{X}^{H} \widetilde{\Sigma}_{RD} \widetilde{X}).$ (57)

Furthermore, according to the left-hand part of the identity (48), we have

$$\det \left(\boldsymbol{I}_{N_{R}} + \sigma_{r}^{-2} \widetilde{\boldsymbol{\Sigma}}_{SR} + \sigma_{r}^{2} \sigma_{d}^{-2} \boldsymbol{X}^{\mathrm{H}} \widetilde{\boldsymbol{\Sigma}}_{RD} \boldsymbol{X} \right) \\ \geq \det \left(\boldsymbol{I}_{N_{R}} + \sigma_{r}^{-2} \widetilde{\boldsymbol{\Sigma}}_{SR} + \sigma_{r}^{2} \sigma_{d}^{-2} \boldsymbol{\Sigma}_{T} \right) \\ = \det \left(\boldsymbol{I}_{N_{R}} + \sigma_{r}^{-2} \widetilde{\boldsymbol{\Sigma}}_{SR} + \sigma_{r}^{2} \sigma_{d}^{-2} \widetilde{\boldsymbol{X}}^{\mathrm{H}} \widetilde{\boldsymbol{\Sigma}}_{RD} \widetilde{\boldsymbol{X}} \right).$$
(58)

The inequality in (58) becomes equality when $X^{H} \tilde{\Sigma}_{RD} X$ is diagonal. By substituting (57) and (58) into (56), we have

$$\frac{2}{B} \operatorname{EE}(\boldsymbol{V}_{S}, \boldsymbol{X}, \boldsymbol{\Delta}_{RD} \otimes \boldsymbol{E}_{RD}) \\
\leq \frac{\log \det \left(\boldsymbol{I}_{N_{R}} + \sigma_{r}^{2} \sigma_{d}^{-2} \widetilde{\boldsymbol{X}}^{\mathrm{H}} \widetilde{\boldsymbol{\Sigma}}_{RD} \widetilde{\boldsymbol{X}}\right)}{\operatorname{Tr}(\boldsymbol{V}_{S}) + \operatorname{Tr}(\sigma_{r}^{2} \widetilde{\boldsymbol{X}} \widetilde{\boldsymbol{X}}^{\mathrm{H}}) + P_{C}} \\
+ \frac{\log \det \left(\boldsymbol{I}_{N_{R}} + \sigma_{r}^{-2} \widetilde{\boldsymbol{\Sigma}}_{SR}\right)}{\operatorname{Tr}(\boldsymbol{V}_{S}) + \operatorname{Tr}(\sigma_{r}^{2} \widetilde{\boldsymbol{X}} \widetilde{\boldsymbol{X}}^{\mathrm{H}}) + P_{C}} \\
- \frac{\log \det \left(\boldsymbol{I}_{N_{R}} + \sigma_{r}^{-2} \widetilde{\boldsymbol{\Sigma}}_{SR} + \sigma_{r}^{2} \sigma_{d}^{-2} \widetilde{\boldsymbol{X}}^{\mathrm{H}} \widetilde{\boldsymbol{\Sigma}}_{RD} \widetilde{\boldsymbol{X}}\right)}{\operatorname{Tr}(\boldsymbol{V}_{S}) + \operatorname{Tr}(\sigma_{r}^{2} \widetilde{\boldsymbol{X}} \widetilde{\boldsymbol{X}}^{\mathrm{H}}) + P_{C}} \\
= \frac{2}{B} \operatorname{EE}(\boldsymbol{V}_{S}, \widetilde{\boldsymbol{X}}, \boldsymbol{\Delta}_{RD} \otimes \boldsymbol{E}_{RD}).$$
(59)

Since the inequality in (59) becomes equality for the diagonal $X^{\mathrm{H}} \widetilde{\Sigma}_{RD} X$, to maximize the EE metric, the optimal X must satisfy $X^{\mathrm{H}} \widetilde{\Sigma}_{RD} X = \Sigma_T$. By introducing the $N_C \times N_C$ diagonal matrix $\overline{\Sigma}_{RD} = \operatorname{diag} \{ \widetilde{\sigma}_{rd,1}^2, \cdots, \widetilde{\sigma}_{rd,N_C}^2 \}$, which is positive definite since the first N_C diagonal elements of $\widetilde{\Sigma}_{RD}$ are positive, and denoting $X^{\mathrm{H}} = [X_1^{\mathrm{H}} X_2^{\mathrm{H}}]$ with $X_1 \in \mathbb{C}^{N_C \times N_R}$ and $X_2 \in \mathbb{C}^{(N_R - N_C) \times N_R}$, we re-express $X^{\mathrm{H}} \widetilde{\Sigma}_{RD} X = \Sigma_T$ as

$$\begin{bmatrix} \boldsymbol{X}_{1}^{\mathrm{H}} \ \boldsymbol{X}_{2}^{\mathrm{H}} \end{bmatrix} \begin{bmatrix} \overline{\boldsymbol{\Sigma}}_{RD} & \boldsymbol{0} \\ \boldsymbol{0} & \boldsymbol{0} \end{bmatrix} \begin{bmatrix} \boldsymbol{X}_{1} \\ \boldsymbol{X}_{2} \end{bmatrix} = \begin{bmatrix} \overline{\boldsymbol{\Sigma}}_{T} & \boldsymbol{0} \\ \boldsymbol{0} & \boldsymbol{0} \end{bmatrix}$$
$$\Rightarrow \boldsymbol{X}_{1}^{\mathrm{H}} \overline{\boldsymbol{\Sigma}}_{RD} \boldsymbol{X}_{1} = \begin{bmatrix} \overline{\boldsymbol{\Sigma}}_{T} & \boldsymbol{0} \\ \boldsymbol{0} & \boldsymbol{0} \end{bmatrix}, \quad (60)$$

where $\overline{\Sigma}_T \in \mathbb{C}^{N_C \times N_C}$ is a positive semidefinite diagonal sub-matrix of Σ_T . (60) indicates that X_2 has no effect on realizing $X^{\mathrm{H}} \widetilde{\Sigma}_{RD} X = \Sigma_T$. Similarly to [6], we can infer from (60) that

$$\boldsymbol{X} = \left[\boldsymbol{X}_{1}^{T}, \boldsymbol{X}_{2}^{T}\right]^{T} = \left[(\overline{\boldsymbol{\Sigma}}_{RD}^{-\frac{1}{2}} \boldsymbol{Q} \widetilde{\boldsymbol{\Sigma}}_{T}^{\frac{1}{2}})^{T}, \boldsymbol{X}_{2}^{T} \right]^{T}, \quad (61)$$

where $\boldsymbol{Q} \in \mathbb{C}^{N_C \times N_C}$ is an arbitrary unitary matrix and $\widetilde{\boldsymbol{\Sigma}}_T^{\frac{1}{2}} = [\overline{\boldsymbol{\Sigma}}_T^{\frac{1}{2}} \ \boldsymbol{0}] \in \mathbb{C}^{N_C \times N_R}$.

To further determine the optimal X_2 , we consider the relay power constraint

$$\operatorname{Tr}(\sigma_r^2 \boldsymbol{X}^{\mathrm{H}} \boldsymbol{X}) = \operatorname{Tr}(\sigma_r^2 \boldsymbol{X}_1 \boldsymbol{X}_1^{\mathrm{H}}) + \operatorname{Tr}(\sigma_r^2 \boldsymbol{X}_2 \boldsymbol{X}_2^{\mathrm{H}})$$
$$\geq \operatorname{Tr}(\sigma_r^2 \boldsymbol{X}_1 \boldsymbol{X}_1^{\mathrm{H}}), \qquad (62)$$

and the last inequality becomes the equality when $X_2 = \mathbf{0}_{(N_R - N_C) \times N_R}$. That is, $X_2 = \mathbf{0}_{(N_R - N_C) \times N_R}$ is the best choice in terms of minimizing relay power consumption. Furthermore,

$$\operatorname{Tr}(\sigma_{r}^{2}\boldsymbol{X}_{1}\boldsymbol{X}_{1}^{\mathrm{H}}) = \operatorname{Tr}\left(\sigma_{r}^{2}\overline{\boldsymbol{\Sigma}}_{RD}^{-\frac{1}{2}}\boldsymbol{Q}\widetilde{\boldsymbol{\Sigma}}_{T}^{\frac{1}{2}}(\widetilde{\boldsymbol{\Sigma}}_{T}^{\frac{1}{2}})^{\mathrm{H}}\boldsymbol{Q}^{\mathrm{H}}(\overline{\boldsymbol{\Sigma}}_{RD}^{-\frac{1}{2}})^{\mathrm{H}}\right) = \operatorname{Tr}(\sigma_{r}^{2}\overline{\boldsymbol{\Sigma}}_{RD}^{-1}\boldsymbol{Q}\overline{\boldsymbol{\Sigma}}_{T}\boldsymbol{Q}^{\mathrm{H}}) \geq \operatorname{Tr}(\sigma_{r}^{2}\overline{\boldsymbol{\Sigma}}_{RD}^{-1}\overline{\boldsymbol{\Sigma}}_{T}),$$

$$(63)$$

where the last inequality is due to the left-hand part of the identity (49) and the equality holds when $Q = I_{N_C}$ according to Lemma 2. Thus the optimal structure of X satisfies

$$\boldsymbol{X}^{\star} = \begin{bmatrix} \boldsymbol{\overline{\Sigma}}_{RD}^{-\frac{1}{2}} \boldsymbol{\widetilde{\Sigma}}_{T}^{\frac{1}{2}} \\ \boldsymbol{0}_{(N_{R}-N_{C})\times N_{R}} \end{bmatrix} = \begin{bmatrix} \boldsymbol{\overline{\Sigma}}_{RD}^{-\frac{1}{2}} \boldsymbol{\overline{\Sigma}}_{T}^{\frac{1}{2}} & \boldsymbol{0} \\ \boldsymbol{0} & \boldsymbol{0} \end{bmatrix}$$
$$= \boldsymbol{\Sigma}_{X} = \operatorname{diag}\{\sigma_{x,1}, \cdots, \sigma_{x,N_{C}}, 0, \cdots, 0\}. \quad (64)$$

Using this optimal X^* in (55), we obtain the optimal structure of the relay beamforming matrix

$$\boldsymbol{W}_{R}^{\star} = \widetilde{\boldsymbol{Q}}_{RD} \boldsymbol{\Sigma}_{X} \left(\boldsymbol{I}_{N_{R}} + \sigma_{r}^{-2} \widetilde{\boldsymbol{\Sigma}}_{SR} \right)^{-\frac{1}{2}} \widetilde{\boldsymbol{U}}_{SR}^{\mathrm{H}}, \quad (65)$$

with

$$\operatorname{EE}(V_{S}, X, \Delta_{RD} \otimes E_{RD}) \leq \operatorname{EE}(V_{S}, \Sigma_{X}, \Delta_{RD} \otimes E_{RD}). \quad (66)$$

2) Optimal V_S^* : Denote the EVD of V_S by $V_S = U_S \Sigma_S U_S^{\text{H}}$, where $U_S \in \mathbb{C}^{N_S \times N_S}$ is the unitary matrix and $\Sigma_S = \text{diag}\{\lambda_{s,1}, \dots, \lambda_{s,N_S}\}$ has the N_S nonnegative diagonal elements. Using the SVD of H_{SR} given in (11), (52) can be rewritten as

$$\boldsymbol{H}_{SR}\boldsymbol{V}_{S}\boldsymbol{H}_{SR}^{\mathrm{H}} = \boldsymbol{U}_{SR}\boldsymbol{\Sigma}_{SR}\boldsymbol{Q}_{SR}^{\mathrm{H}}\boldsymbol{U}_{S}\boldsymbol{\Sigma}_{S}\boldsymbol{U}_{S}^{\mathrm{H}}\boldsymbol{Q}_{SR}\boldsymbol{\Sigma}_{SR}^{\mathrm{H}}\boldsymbol{U}_{SR}^{\mathrm{H}}$$
$$= \widetilde{\boldsymbol{U}}_{SR}\widetilde{\boldsymbol{\Sigma}}_{SR}\widetilde{\boldsymbol{U}}_{SR}^{\mathrm{H}}.$$
(67)

Furthermore, the EE metric in the problem (10) can be re-expressed as

$$=\frac{\log \det(\boldsymbol{I}_{N_{S}}, \boldsymbol{W}_{R}, \boldsymbol{\Delta}_{RD} \otimes \boldsymbol{E}_{RD})}{\operatorname{Tr}(\boldsymbol{V}_{S}) + \operatorname{Tr}(\boldsymbol{W}_{R}(\boldsymbol{H}_{SR}\boldsymbol{V}_{S}\boldsymbol{H}_{RH}^{\mathrm{H}}\boldsymbol{H}_{SR}^{\mathrm{H}}\boldsymbol{V}_{S}\boldsymbol{H}_{SR}^{\mathrm{H}} + \sigma_{r}^{2}\boldsymbol{I}_{N_{R}})\boldsymbol{W}_{R}^{\mathrm{H}}) + P_{C}},$$
(68)

where $\widetilde{\boldsymbol{V}}_{S} = \boldsymbol{U}_{S}\boldsymbol{\Sigma}_{S}^{\frac{1}{2}}$ and $\boldsymbol{B} = \sigma_{r}^{2}\boldsymbol{H}_{RD}\boldsymbol{W}_{R}\boldsymbol{W}_{R}^{\mathrm{H}}\boldsymbol{H}_{RD}^{\mathrm{H}} + \sigma_{d}^{2}\boldsymbol{I}_{N_{D}}$.

Substituting the optimal \boldsymbol{W}_R^{\star} of (65) into (68) yields

$$\frac{\frac{2}{B} \cdot \text{EE}\left(\boldsymbol{V}_{S}, \boldsymbol{\Sigma}_{X}, \boldsymbol{\Delta}_{RD} \otimes \boldsymbol{E}_{RD}\right)}{\log \det\left(\boldsymbol{I}_{N_{S}} + \widetilde{\boldsymbol{V}}_{S}^{\text{H}} \boldsymbol{H}_{SR}^{\text{H}} \widetilde{\boldsymbol{U}}_{SR} \widetilde{\boldsymbol{\Sigma}}_{RDX} \widetilde{\boldsymbol{U}}_{SR}^{\text{H}} \boldsymbol{H}_{SR} \widetilde{\boldsymbol{V}}_{S}\right)}{\operatorname{Tr}(\boldsymbol{\Sigma}_{S}) + \operatorname{Tr}\left(\sigma_{r}^{2} \boldsymbol{\Sigma}_{X} \boldsymbol{\Sigma}_{X}\right) + P_{C}},$$
(69)

with

0

$$\widetilde{\boldsymbol{\Sigma}}_{RDX} = \boldsymbol{\Sigma}_{RDX}^{\mathrm{H}} \left(\sigma_d^2 \boldsymbol{I}_{N_R} + \sigma_r^2 \boldsymbol{\Sigma}_{RDX} \boldsymbol{\Sigma}_{RDX}^{\mathrm{H}} \right)^{-1} \boldsymbol{\Sigma}_{RDX}, \quad (70)$$

where $\Sigma_{RDX} = \widetilde{\Sigma}_{RD}^{\frac{1}{2}} \Sigma_X (I_{N_R} + \sigma_r^{-2} \widetilde{\Sigma}_{SR})^{-\frac{1}{2}}$. Since rank (Σ_X) = rank $(\widetilde{\Sigma}_{RD}) = N_C$, the diagonal matrix $\widetilde{\Sigma}_{RDX}$ also has the rank of N_C . According to (67), we can re-express (69) as

$$\frac{\frac{2}{B} \cdot \text{EE}(\boldsymbol{V}_{S}, \boldsymbol{\Sigma}_{X}, \boldsymbol{\Delta}_{RD} \otimes \boldsymbol{E}_{RD})}{\log \det \left(\boldsymbol{I}_{N_{R}} + \overline{\boldsymbol{U}}_{SR}^{H} \boldsymbol{\Sigma}_{SR} \overline{\boldsymbol{Q}}_{SR} \boldsymbol{\Sigma}_{S} \overline{\boldsymbol{Q}}_{SR}^{H} \boldsymbol{\Sigma}_{SR}^{H} \overline{\boldsymbol{U}}_{SR} \widetilde{\boldsymbol{\Sigma}}_{RDX}\right)}{\operatorname{Tr}(\boldsymbol{\Sigma}_{S}) + \operatorname{Tr}(\sigma_{r}^{2} \boldsymbol{\Sigma}_{X} \boldsymbol{\Sigma}_{X}) + P_{C}} \\
\leq \frac{\log \det \left(\boldsymbol{I}_{N_{R}} + \boldsymbol{\Sigma}_{SR} \boldsymbol{\Sigma}_{S} \boldsymbol{\Sigma}_{SR}^{H} \widetilde{\boldsymbol{\Sigma}}_{RDX}\right)}{\operatorname{Tr}(\boldsymbol{\Sigma}_{S}) + \operatorname{Tr}(\sigma_{r}^{2} \boldsymbol{\Sigma}_{X} \boldsymbol{\Sigma}_{X}) + P_{C}} \\
= \frac{2}{B} \cdot \operatorname{EE}(\boldsymbol{\Sigma}_{S}, \boldsymbol{\Sigma}_{X}, \boldsymbol{\Delta}_{RD} \otimes \boldsymbol{E}_{RD}), \quad (71)$$

where $\overline{U}_{SR}=U_{SR}^{\mathrm{H}}\widetilde{U}_{SR}$ and $\overline{Q}_{SR}=Q_{SR}^{\mathrm{H}}U_{S}$ are unitary matrices. According to the right-hand part of identity (48), it can be inferred that det $(I_{N}+AB) \leq \prod_{i=1}^{N} (1+\lambda_{i}(A)\lambda_{i}(B))$ and the inequality in (71) becomes equality when $\overline{U}_{SR}=I_{N_{R}}$ and $\overline{Q}_{SR}=I_{N_{S}}$. The diagonal matrix $\widetilde{\Sigma}_{RDX}$ (70) is satisfied at the maximum EE point. As a result, the optimal structures of \widetilde{U}_{SR} and U_{S} are $\widetilde{U}_{SR}=U_{SR}$ and $U_{S}=Q_{SR}$. Accordingly, the optimal structure of V_{S} is given by

$$\boldsymbol{V}_{S}^{\star} = \boldsymbol{Q}_{SR} \boldsymbol{\Sigma}_{S} \boldsymbol{Q}_{SR}^{\mathrm{H}}.$$
 (72)

Combining (67) with (72) leads to $\boldsymbol{H}_{SR}\boldsymbol{V}_{S}^{\star}\boldsymbol{H}_{SR}^{\mathrm{H}} = \boldsymbol{U}_{SR}\boldsymbol{\Sigma}_{SR} \cdot \boldsymbol{\Sigma}_{S}\boldsymbol{\Sigma}_{SR}\boldsymbol{U}_{SR}^{\mathrm{H}} = \boldsymbol{\widetilde{U}}_{SR}\boldsymbol{\widetilde{\Sigma}}_{SR}\boldsymbol{\widetilde{U}}_{SR}^{\mathrm{H}}$ and due to $\boldsymbol{\widetilde{U}}_{SR} = \boldsymbol{U}_{SR}$,

we also have $\widetilde{\Sigma}_{SR} = \Sigma_{SR} \Sigma_S \Sigma_{SR}^H$. Thus the optimal structure of the relay beamforming matrix in (65) can be re-expressed as

$$\boldsymbol{W}_{R}^{\star} = \widetilde{\boldsymbol{Q}}_{RD} \boldsymbol{\Sigma}_{X} \left(\boldsymbol{I}_{N_{R}} + \sigma_{r}^{-2} \boldsymbol{\Sigma}_{SR} \boldsymbol{\Sigma}_{S} \boldsymbol{\Sigma}_{SR}^{\mathrm{H}} \right)^{-\frac{1}{2}} \boldsymbol{U}_{SR}^{\mathrm{H}}.$$
 (73)

3) Worst-Case $\Delta_{RD}^* \otimes E_{RD}^*$: First, we introduce the following lemma [15], [37].

Lemma 3: For $\mathbf{A} \in \mathbb{C}^{N_u \times N_l}$, $\mathbf{B}, \mathbf{C} \in \mathbb{C}^{N_l \times N_m}$ with rank $(\mathbf{A}) = \operatorname{rank}(\mathbf{B}) = \operatorname{rank}(\mathbf{C}) = N = \min\{N_u, N_l, N_m\}$, whose SVs are $\sigma_i(\mathbf{A})$, $\sigma_i(\mathbf{B})$ and $\sigma_i(\mathbf{C})$, $i = 1, \dots, N$, respectively, we have

$$\sigma_N(\boldsymbol{A})\sigma_i(\boldsymbol{B}) \le \sigma_i(\boldsymbol{A}\boldsymbol{B}),\tag{74}$$

$$\left(\sigma_i(\boldsymbol{B}) - \sigma_1(\boldsymbol{C})\right)^+ \le \sigma_i(\boldsymbol{B} + \boldsymbol{C}). \tag{75}$$

(74) and (75) become equalities only if $\{A, B\}$ and $\{B, C\}$ are simultaneously diagonalizable.

Based on the definitions of the diagonal matrices Σ_S , Σ_X and $\widetilde{\Sigma}_{RDX}$ as well as the rectangular diagonal matrix Σ_{SR} , at V_S^* and W_R^* , the achievable EE metric in (71) can be expressed as (76) at the top of the next page, where $N_L = \min \{N_P, N_C\}$. It is readily observed from the first term of the numerator in (76) that the minimum value of $\text{EE}(\Sigma_S, \Sigma_X, \Delta_{RD} \otimes E_{RD})$ is attained only when every $\widetilde{\sigma}_{rd,i}^2$ realizes its minimum, subject to the spectral norm constraint. Note that $\widetilde{\sigma}_{rd,i}^2$, $1 \le i \le N_C$, are unknown since they are related to the unknown CSI H_{RD} . However, for the known nominal CSI \widehat{H}_{RD} , we have the SVD $\widehat{H}_{RD} = \widehat{U}_{RD}\widehat{\Sigma}_{RD}\widehat{Q}_{RD}^{H}$, in which the $N_D \times N_R$ diagonal rectangular matrix $\widehat{\Sigma}_{RD}$ contains N_C positive elements $\{\widehat{\sigma}_{rd,1}, \dots, \widehat{\sigma}_{rd,N_C}\}$.

3.1) Additive CSI errors: $\boldsymbol{H}_{RD}^{H}\boldsymbol{H}_{RD} = \left(\widehat{\boldsymbol{H}}_{RD} + \boldsymbol{\Delta}_{RD}\right)^{H} \cdot \left(\widehat{\boldsymbol{H}}_{RD} + \boldsymbol{\Delta}_{RD}\right) = \widetilde{\boldsymbol{Q}}_{RD}\widetilde{\boldsymbol{\Sigma}}_{RD}\widetilde{\boldsymbol{Q}}_{RD}^{H}$. Applying Lemma 3 and considering $\left\|\boldsymbol{\Delta}_{RD}\right\|_{2} = \sigma_{1}\left(\boldsymbol{\Delta}_{RD}\right) \leq \epsilon_{a}$, we conclude

$$\widetilde{\sigma}_{rd,i} \ge \left(\widehat{\sigma}_{rd,i} - \epsilon_a\right)^+, \quad 1 \le i \le N_C.$$
(77)

All the inequalities in (77) become the equalities when $\widetilde{Q}_{RD} = \widehat{Q}_{RD}$ according to Lemma 3, and the minimum $\text{EE}(\Sigma_S, \Sigma_X, \Delta_{RD})$ is attained with $\widetilde{\sigma}_{rd,i} = (\widehat{\sigma}_{rd,i} - \epsilon_a)^+$, $1 \le i \le N_C$. Consequently, the resultant worst-case CSI error is given by

$$\boldsymbol{\Delta}_{RD}^{\star} = -\widehat{\boldsymbol{U}}_{RD} \begin{bmatrix} \widetilde{\boldsymbol{\Lambda}}_{RD} & \boldsymbol{0}_{N_{C} \times (N_{R} - N_{C})} \\ \boldsymbol{0}_{(N_{D} - N_{C}) \times N_{C}} & \boldsymbol{0}_{(N_{D} - N_{C}) \times (N_{R} - N_{C})} \end{bmatrix} \widehat{\boldsymbol{Q}}_{RD}^{\mathrm{H}}$$

$$= -\widehat{\boldsymbol{U}}_{RD} \boldsymbol{\Lambda}_{RD} \widehat{\boldsymbol{Q}}_{RD}^{\mathrm{H}},$$

$$(78)$$

where $\widetilde{\Lambda}_{RD} = \text{diag} \{\min \{\widehat{\sigma}_{rd,1}, \epsilon_a\}, \cdots, \min \{\widehat{\sigma}_{rd,N_C}, \epsilon_a\}\},\$ and we also have

$$\operatorname{EE}(\boldsymbol{\Sigma}_{S},\boldsymbol{\Sigma}_{X},\boldsymbol{\Delta}_{RD}) \geq \operatorname{EE}(\boldsymbol{\Sigma}_{S},\boldsymbol{\Sigma}_{X},\boldsymbol{\Delta}_{RD}^{\star}). \tag{79}$$

Observe that \widehat{U}_{RD} , Λ_{RD} and $\widehat{Q}_{RD}^{\text{H}}$ in (78) are all known.

3.2) Multiplicative CSI errors: $\boldsymbol{H}_{RD}^{\mathrm{H}} \boldsymbol{H}_{RD} = \widehat{\boldsymbol{H}}_{RD}^{\mathrm{H}} (\boldsymbol{I}_{N_D} + \boldsymbol{E}_{RD})^{\mathrm{H}} (\boldsymbol{I}_{N_D} + \boldsymbol{E}_{RD}) \widehat{\boldsymbol{H}}_{RD} = \widetilde{\boldsymbol{Q}}_{RD} \widetilde{\boldsymbol{\Sigma}}_{RD} \widetilde{\boldsymbol{Q}}_{RD}^{\mathrm{H}}$. Similarly to the case of additive CSI errors, based on Lemma 3 and $\|\boldsymbol{E}_{RD}\|_2 = \sigma_1(\boldsymbol{E}_{RD}) \leq \epsilon_m$, the minimum values of $\widetilde{\sigma}_{rd,i}, \forall i$, required for minimizing EE $(\boldsymbol{\Sigma}_S, \boldsymbol{\Sigma}_X, \boldsymbol{E}_{RD})$ are obtained as

$$\widetilde{\sigma}_{rd,i} \ge \sigma_{N_D} \left(\boldsymbol{I}_{N_D} + \boldsymbol{E}_{RD} \right) \widehat{\sigma}_{rd,i} \ge (1 - \epsilon_m)^+ \widehat{\sigma}_{rd,i}, \\ 1 < i < N_C. \quad (80)$$

0

$$\frac{\frac{2}{B} \cdot \text{EE}\left(\boldsymbol{\Sigma}_{S}, \boldsymbol{\Sigma}_{X}, \boldsymbol{\Delta}_{RD} \otimes \boldsymbol{E}_{RD}\right) = \frac{\sum_{i=1}^{N_{L}} \log\left(\frac{1 + \sigma_{r}^{2} \sigma_{d}^{-2} \sigma_{x,i}^{2} \tilde{\sigma}_{r,i}^{2}}{1 + \sigma_{r}^{-2} \sigma_{sr,i}^{2} \lambda_{s,i} + \sigma_{r}^{2} \sigma_{d}^{-2} \sigma_{x,i}^{2} \tilde{\sigma}_{r,i}^{2}}\right) + \sum_{i=1}^{N_{L}} \log\left(1 + \sigma_{r}^{-2} \sigma_{sr,i}^{2} \lambda_{s,i}\right)}{\sum_{i=1}^{N_{S}} \lambda_{s,i} + \sum_{i=1}^{N_{C}} \sigma_{r}^{2} \sigma_{x,i}^{2}} + P_{C}} = \frac{\sum_{i=1}^{N_{L}} \log\left(1 - \frac{\sigma_{r}^{-2} \sigma_{sr,i}^{2} \lambda_{s,i}}{1 + \sigma_{r}^{-2} \sigma_{sr,i}^{2} \lambda_{s,i} + \sigma_{r}^{2} \sigma_{\sigma}^{-2} \sigma_{x,i}^{2}}\right) + \sum_{i=1}^{N_{L}} \log\left(1 + \sigma_{r}^{-2} \sigma_{sr,i}^{2} \lambda_{s,i}\right)}{\sum_{i=1}^{N_{S}} \lambda_{s,i} + \sum_{i=1}^{N_{C}} \sigma_{r}^{2} \sigma_{x,i}^{2}} + P_{C}}, \quad (76)$$

The two inequalities in (80) simultaneously become the equalities when $\tilde{Q}_{RD} = \hat{Q}_{RD}$ according to Lemma 3. Therefore, the resulting worst-case E_{RD}^{\star} is given by $E_{RD}^{\star} = -\epsilon_m I_{N_D}$, which only depends on the known ϵ_m , and we naturally have

$$\operatorname{EE}\left(\boldsymbol{\Sigma}_{S}, \boldsymbol{\Sigma}_{X}, \boldsymbol{E}_{RD}\right) \geq \operatorname{EE}\left(\boldsymbol{\Sigma}_{S}, \boldsymbol{\Sigma}_{X}, \boldsymbol{E}_{RD}^{\star}\right). \tag{81}$$

Because $\tilde{Q}_{RD} = \hat{Q}_{RD}$ holds for both additive and multiplicative CSI errors, (73) becomes

$$\boldsymbol{W}_{R}^{\star} = \widehat{\boldsymbol{Q}}_{RD} \boldsymbol{\Sigma}_{X} \left(\boldsymbol{I}_{N_{R}} + \sigma_{r}^{-2} \boldsymbol{\Sigma}_{SR} \boldsymbol{\Sigma}_{S} \boldsymbol{\Sigma}_{SR}^{\mathrm{H}} \right)^{-\frac{1}{2}} \boldsymbol{U}_{SR}^{\mathrm{H}}.$$
(82)

Observe that \widehat{Q}_{RD} , Σ_{SR} and U_{SR} are all known, while Σ_S and Σ_X are the new optimization variables. This completes the proof.

B. Proof of Theorem 2

Proof: From (66) and (71) in Appendix A, it is easily seen that $EE(V_S^*, W_R^*, \Delta_{RD}^* \otimes E_{RD}^*) \geq EE(V_S, W_R, \Delta_{RD}^* \otimes E_{RD}^*)$ holds for any feasible V_S and W_R . Similarly, from (79) and (81) in Appendix A, it is seen that $EE(V_S^*, W_R^*, \Delta_{RD}^* \otimes E_{RD}^*) \leq EE(V_S^*, W_R^*, \Delta_{RD} \otimes E_{RD})$ holds for any feasible $\Delta_{RD} \otimes E_{RD}$. Thus the optimal $\{V_S^*, W_R^*, \Delta_{RD}^* \otimes E_{RD}^*\}$ is a saddle point of the original robust EE optimization problem (8). This completes the proof.

C. Proof of Theorem 3

Proof: According to [11], for the statistically imperfect source-relay channel H_{SR} , the channel-diagonalizing structure is optimal for the source covariance matrix V_S^{\star} and the relay beamforming W_{R}^{\star} for any relay-destination channel H_{RD} . That is, the eigenvectors of the optimal V_{S}^{\star} are aligned with that of the source correlation matrix R_S , while the left and right singular matrices of the optimal W_R^{\star} are aligned with the right singular matrix of the relay-destination channel H_{RD} and the eigenvectors of the relay correlation matrix R_R , respectively [11]. Based on the optimal V_S^{\star} and W_R^{\star} with the channel-diagonalizing structure (29), we naturally obtain the same worst-case error ${f \Delta}^{\star}_{RD} arnothing E^{\star}_{RD}$ as that given in Theorem 1 by utilizing Lemma 3 of Appendix A. Moreover, we can also conclude that the solutions provided in Theorem 3 are the saddle point of $EE(V_S, W_R, \Delta_{RD} \otimes E_{RD})$ by referring to the proof of Theorem 2.

D. Proof of Theorem 4

Proof: Construct the min-max EE counterpart problem to the problem (33). First by fixing the relay beamforming matrices of relays $k \in \{2, \dots, K\}$ and referring to the proof of Theorem 1, we obtain the optimal beamforming matrix of relay k' = 1 and the optimal source covariance matrix, both having channel-diagonalizing structure, as well as obtain the same worst-case CSI error as given in Theorem 1. In a similar manner, the remaining optimal beamforming matrices of relays $k', 2 \le k' \le K$, with channel-diagonalizing structure can be obtained one by one by fixing the relay beamforming matrices of relays $k \in \{1, 2, \dots, K\} \setminus k'$ and given the optimal source covariance matrix. This proves that the solution of (42) to (44) form the optimal solution of this min-max EE counterpart problem. Then referring to the proof of Theorem 2, we conclude that the solution of (42) to (44) is a saddle point of $\operatorname{EE}_M(V_S, \widetilde{W}_R, \Delta_{R_K D} \otimes \boldsymbol{E}_{R_K D}).$

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